

MAIT 627 Fast Multipole Methods

Lecture 7

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Outline

- Norm of the translation operator
- Example of SIR-translation
- Summary of requirements for functions (potentials) that can be used in FMM
- Idea of a Single Level FMM (SLFMM)
- Space division and expansion domains
- SLFMM algorithm
- Asymptotic complexity of SLFMM
- Optimization of SLFMM

Example from previous lectures

$$\Phi(y, x_i) = \frac{1}{y - x_i}$$

$|y - x_*| < |x_i - x_*| :$

R-expansion

$$\Phi(y, x_i) = \sum_{m=0}^{\infty} a_m(x_i, x_*) R_m(y - x_*),$$

$$a_m(x_i, x_*) = -(x_i - x_*)^{-m-1}, \quad m = 0, 1, \dots,$$

$$R_m(y - x_*) = (y - x_*)^m, \quad m = 0, 1, \dots$$

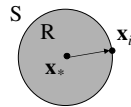
$|y - x_*| > |x_i - x_*| :$

S-expansion

$$\Phi(y, x_i) = \sum_{m=0}^{\infty} b_m(x_i, x_*) S_m(y - x_*),$$

$$b_m(x_i, x_*) = (x_i - x_*)^m, \quad m = 0, 1, \dots,$$

$$S_m(y - x_*) = (y - x_*)^{-m-1}, \quad m = 0, 1, \dots$$



In this case we have

$(|y - x_*| < |t|)$

$$S_n(y - x_* + t) = (t + y)^{-n-1} = \sum_{m=0}^{\infty} \frac{1}{m!} \frac{d^m S_n(t)}{dt^m} (y - x_*)^m$$

$$= \sum_{m=0}^{\infty} \frac{1}{m!} \frac{d^m S_n(t)}{dt^m} R_m(y - x_*) = \sum_{m=0}^{\infty} (S|R)_{mn}(t) R_m(y - x_*).$$

So

$$(S|R)_{mn}(t) = \frac{1}{m!} \frac{d^m S_n(t)}{dt^m} = \frac{(-1)^m (m+n)!}{m! n! t^{m+n+1}}.$$

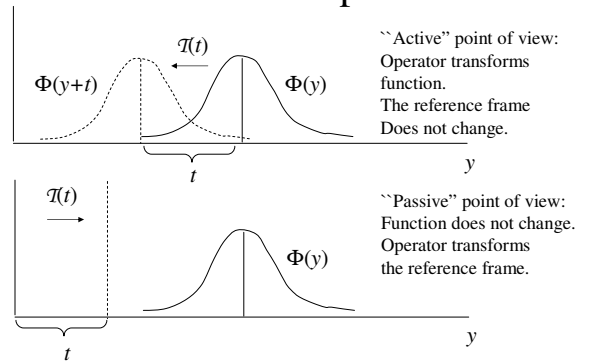
$$(S|R)(t) = \begin{pmatrix} t^{-1} & t^{-2} & t^{-3} & \dots \\ -t^{-2} & -2t^{-3} & -3t^{-4} & \dots \\ t^{-3} & 3t^{-4} & 6t^{-5} & \dots \\ \dots & \dots & \dots & \dots \end{pmatrix}$$

Norm of the Translation Operator

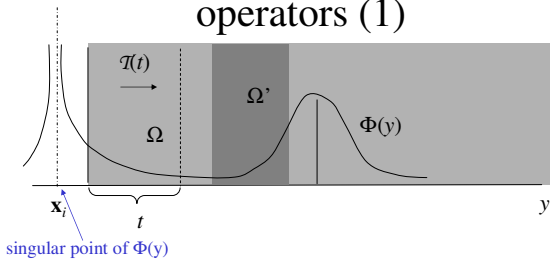
Theorem. Let $\mathbb{F}(\Omega)$ be a set of functions bounded in \mathbb{R}^d . Then $\|\mathcal{T}(t)\| = 1$.
Proof.

$$\|\mathcal{T}(t)\| = \frac{\|\mathcal{T}(t)\Phi(y)\|}{\|\Phi(y)\|} = \frac{\|\Phi(y+t)\|}{\|\Phi(y)\|} = \frac{\sup_{y \in \mathbb{R}^d} |\Phi(y+t)|}{\sup_{y \in \mathbb{R}^d} |\Phi(y)|} = 1.$$

Active and Passive points of view on translation operator



Norms of RIR, SIS, and SIR-operators (1)



$\Phi(y)$ is bounded in Ω .
 $\Omega' \subset \Omega$.

Therefore $\Phi(y)$ is bounded in Ω' , and

$$\|\Phi(y)\|_{\Omega'} = \sup_{y \in \Omega'} |\Phi(y)| \leq \sup_{y \in \Omega} |\Phi(y)| = \|\Phi(y)\|_{\Omega}$$

Norms of RIR, SIS, and SIR-operators (2)

From the passive point of view, the translation operator does nothing, but just changes the reference frame. So if we consider that RIR, SIS, and SIR do just change of the reference frame **PLUS they shrink the domain, where the function is bounded, then their norms do not exceed 1.**

$$\Omega' \subset \Omega$$

$$\|(\mathcal{R}|\mathcal{R})(t)\| = \frac{\sup_{y \in \Omega'} |\Phi(y)|}{\sup_{y \in \Omega} |\Phi(y)|} \leq 1,$$

$$\|(\mathcal{S}|\mathcal{S})(t)\| = \frac{\sup_{y \in \Omega'} |\Phi(y)|}{\sup_{y \in \Omega} |\Phi(y)|} \leq 1,$$

$$\|(\mathcal{S}|\mathcal{R})(t)\| = \frac{\sup_{y \in \Omega'} |\Phi(y)|}{\sup_{y \in \Omega} |\Phi(y)|} \leq 1.$$

This is the difference between general translation operator and RIR, SIS, and SIR operators.

Error of exact $R|R$, $S|S$, and $S|R$ -translation

If

$$\|\Phi(\mathbf{y}) - \Phi^p(\mathbf{y})\| < \epsilon,$$

then

$$\begin{aligned} \|(\mathcal{R}|\mathcal{R})(\Phi(\mathbf{y}) - \Phi^p(\mathbf{y}))\| &= \|(\mathcal{R}|\mathcal{R})(\mathbf{1})\| \|\Phi(\mathbf{y}) - \Phi^p(\mathbf{y})\| < \epsilon, \\ \|(\mathcal{S}|\mathcal{S})(\Phi(\mathbf{y}) - \Phi^p(\mathbf{y}))\| &= \|(\mathcal{S}|\mathcal{S})(\mathbf{1})\| \|\Phi(\mathbf{y}) - \Phi^p(\mathbf{y})\| < \epsilon, \\ \|(\mathcal{S}|\mathcal{R})(\Phi(\mathbf{y}) - \Phi^p(\mathbf{y}))\| &= \|(\mathcal{S}|\mathcal{R})(\mathbf{1})\| \|\Phi(\mathbf{y}) - \Phi^p(\mathbf{y})\| < \epsilon. \end{aligned}$$

Four Key Stones of FMM

- Factorization
- Error
- Translation
- Grouping

Summary of formal requirements for functions that can be used in FMM

- We have two sets of points:

$$\mathbf{X} = \{\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_N\}, \quad \mathbf{x}_i \in \mathbb{R}^d, \quad i = 1, \dots, N,$$

$$\mathbf{Y} = \{\mathbf{y}_1, \mathbf{y}_2, \dots, \mathbf{y}_M\}, \quad \mathbf{y}_j \in \mathbb{R}^d, \quad j = 1, \dots, M.$$

- We have functions (potentials):

$$\Phi(\mathbf{x}_i, \mathbf{y}) : \mathbb{R}^d \rightarrow \mathbb{R}, \quad \mathbf{y} \in \mathbb{R}^d, \quad i = 1, \dots, N.$$

- These functions can be factorized as (local expansion):

$$\Phi(\mathbf{x}_i, \mathbf{y}) = \mathbf{A}(\mathbf{x}_i, \mathbf{x}_*) \circ \mathbf{R}(\mathbf{y} - \mathbf{x}_*), \quad |\mathbf{y} - \mathbf{x}_*| < r < |\mathbf{x}_i - \mathbf{x}_*|, \quad i = 1, \dots, N$$

- These functions can be factorized as (far field expansion):

$$\Phi(\mathbf{x}_i, \mathbf{y}) = \mathbf{B}(\mathbf{x}_i, \mathbf{x}_*) \circ \mathbf{S}(\mathbf{x} - \mathbf{x}_*), \quad |\mathbf{y} - \mathbf{x}_*| > R > |\mathbf{x}_i - \mathbf{x}_*|, \quad i = 1, \dots, N$$

- The product is distributive operation with respect to addition

$$(u_1 \mathbf{A}_1 + u_2 \mathbf{A}_2) \circ \mathbf{F} = u_1 \mathbf{A}_1 \circ \mathbf{F} + u_2 \mathbf{A}_2 \circ \mathbf{F}, \quad \mathbf{F} = \mathbf{S}, \mathbf{R}$$

Summary of formal requirements for functions that can be used in FMM (2)

- R -expansion coefficients can be $R|R$ -translated:

$$|\mathbf{x} - \mathbf{x}_{*2}| < |\mathbf{x}_i - \mathbf{x}_{*1}| - |\mathbf{x}_{*1} - \mathbf{x}_{*2}| :$$

$$\mathbf{A}(\mathbf{x}_i, \mathbf{x}_{*2}) = (\mathbf{R}|\mathbf{R})(\mathbf{x}_{*2} - \mathbf{x}_{*1}) \mathbf{A}(\mathbf{x}_i, \mathbf{x}_{*1})$$

- S -expansion coefficients can be $S|S$ -translated:

$$|\mathbf{x} - \mathbf{x}_{*2}| > |\mathbf{x}_{*1} - \mathbf{x}_{*2}| + |\mathbf{x}_i - \mathbf{x}_{*1}|,$$

$$\mathbf{B}(\mathbf{x}_i, \mathbf{x}_{*2}) = (\mathbf{S}|\mathbf{S})(\mathbf{x}_{*2} - \mathbf{x}_{*1}) \mathbf{B}(\mathbf{x}_i, \mathbf{x}_{*1})$$

- S -expansion coefficients can be $S|R$ -translated (converted to R -expansion coefficients)

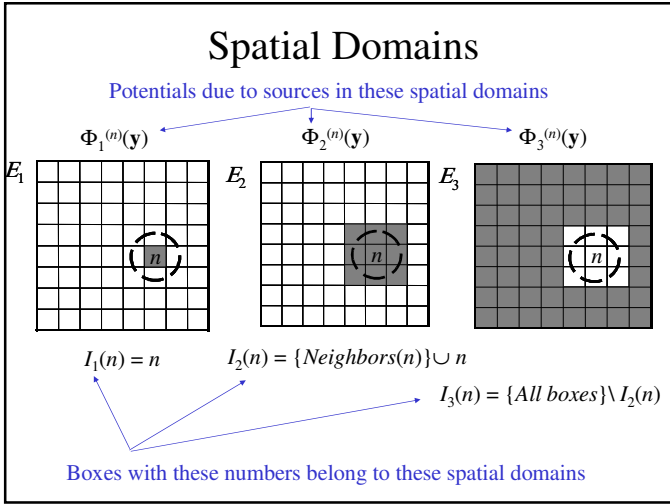
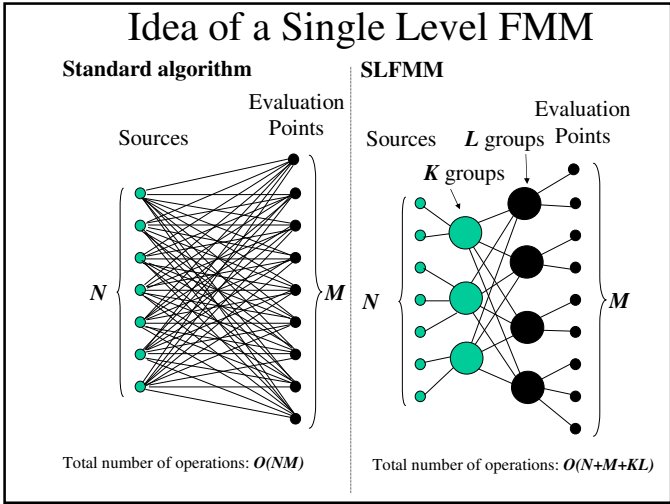
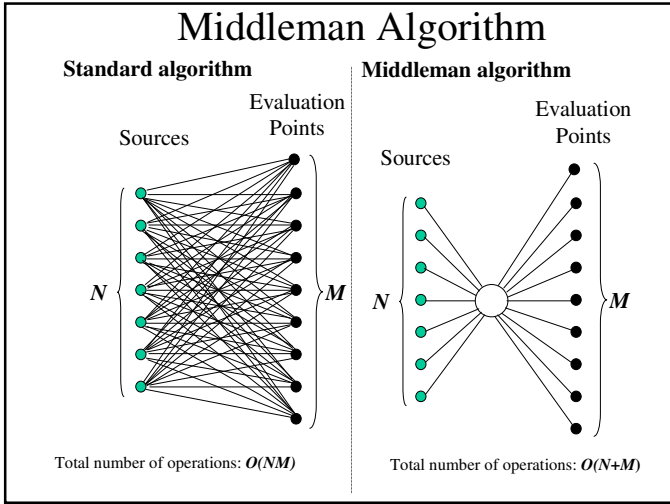
$$|\mathbf{x} - \mathbf{x}_{*2}| < |\mathbf{x}_{*1} - \mathbf{x}_{*2}| + |\mathbf{x}_i - \mathbf{x}_{*1}|,$$

$$\mathbf{A}(\mathbf{x}_i, \mathbf{x}_{*2}) = (\mathbf{S}|\mathbf{R})(\mathbf{x}_{*2} - \mathbf{x}_{*1}) \mathbf{B}(\mathbf{x}_i, \mathbf{x}_{*1})$$

- And we are looking for sums:

$$\mathbf{v}_j = \sum_{i=1}^N u_i \Phi(\mathbf{y}_j, \mathbf{x}_i), \quad j = 1, \dots, M.$$

- Some generalization are possible, say instead of $\Phi(\mathbf{y}_j, \mathbf{x}_i)$ we can consider $\Phi_j(\mathbf{y}_j)$, etc.



Definition of potentials

$$\Phi_1^{(n)}(\mathbf{y}) = \sum_{\mathbf{x}_i \in E_1(n)} u_i \Phi(\mathbf{y}, \mathbf{x}_i),$$

$$\Phi_2^{(n)}(\mathbf{y}) = \sum_{\mathbf{x}_i \in E_2(n)} u_i \Phi(\mathbf{y}, \mathbf{x}_i),$$

$$\Phi_3^{(n)}(\mathbf{y}) = \sum_{\mathbf{x}_i \in E_3(n)} u_i \Phi(\mathbf{y}, \mathbf{x}_i),$$

Since domains $E_2(n)$ and $E_3(n)$ are complementary:

$$\Phi(\mathbf{y}) = \sum_{i=1}^N u_i \Phi(\mathbf{y}, \mathbf{x}_i) = \sum_{\mathbf{x}_i \in E_2(n) \cup E_3(n)} u_i \Phi(\mathbf{y}, \mathbf{x}_i) = \Phi_2^{(n)}(\mathbf{y}) + \Phi_3^{(n)}(\mathbf{y})$$

for arbitrary n .

SLFMM Algorithm

Step 1. Generate S-expansion coefficients
for each box

$$\Phi_1^{(n)}(\mathbf{x}) = \mathbf{C}^{(n)} \circ \mathbf{S}(\mathbf{x} - \mathbf{x}_c^{(n)}),$$

$$\mathbf{C}^{(n)} = \sum_{\mathbf{x}_i \in E_1(n)} u_i \mathbf{B}(\mathbf{x}_i, \mathbf{x}_c^{(n)}).$$

loop over all non-empty source boxes

For $n \in \text{NonEmptySource}$

Get $\mathbf{x}_c^{(n)}$, the center of the box;

$\mathbf{C}^{(n)} = \mathbf{0}$;

For $\mathbf{x}_i \in E_1(n)$ ← loop over all sources in the box

Get $\mathbf{B}(\mathbf{x}_i, \mathbf{x}_c^{(n)})$, the S-expansion coefficients
near the center of the box;

$\mathbf{C}^{(n)} = \mathbf{C}^{(n)} + u_i \mathbf{B}(\mathbf{x}_i, \mathbf{x}_c^{(n)})$;

End;

End;

Implementation can be different!
All we need is to get $\mathbf{C}^{(n)}$.

SLFMM Algorithm

Step 2. (SIR)-translate expansion coefficients

$$\Phi_3^{(n)}(\mathbf{y}) = \mathbf{D}^{(n)} \circ \mathbf{R}(\mathbf{y} - \mathbf{x}_c^{(n)}),$$

$$\mathbf{D}^{(n)} = \sum_{m \in I_3(n)} (\text{SIR})(\mathbf{x}_c^{(n)} - \mathbf{x}_c^{(m)}) \mathbf{C}^{(m)}.$$

loop over all non-empty
evaluation boxes

For $n \in \text{NonEmptyEvaluation}$

Get $\mathbf{x}_c^{(n)}$, the center of the box;

$\mathbf{D}^{(n)} = \mathbf{0}$;

For $m \in I_3(n)$ ← loop over all non-empty source boxes
outside the neighborhood of the n -th box

Get $\mathbf{x}_c^{(m)}$, the center of the box;

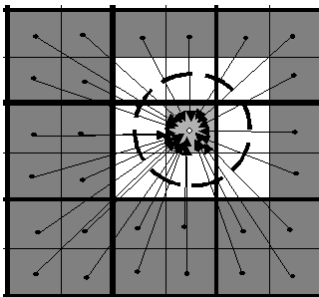
$\mathbf{D}^{(n)} = \mathbf{D}^{(n)} + (\text{SIR})(\mathbf{x}_c^{(n)} - \mathbf{x}_c^{(m)}) \mathbf{C}^{(m)}$;

End;

End;

Implementation can be different!
All we need is to get $\mathbf{D}^{(n)}$.

SIR-translation



SLFMM Algorithm

Step 3. Final Summation

$$v_j = \Phi(y_j) = \sum_{\mathbf{x}_i \in E_2(n)} \Phi(y_j, \mathbf{x}_i) + \mathbf{D}^{(n)} \circ \mathbf{R}(y_j - \mathbf{x}_c^{(n)}), \quad \mathbf{y}_j \in E_1(n).$$

For $n \in \text{NonEmptyEvaluation}$ ← loop over all boxes
containing evaluation points

Get $\mathbf{x}_c^{(n)}$, the center of the box;

For $\mathbf{y}_j \in E_1(n)$ ← loop over all evaluation points in the box

$v_j = \mathbf{D}^{(n)} \circ \mathbf{R}(\mathbf{y}_j - \mathbf{x}_c^{(n)})$;

For $\mathbf{x}_i \in E_2(n)$ ← loop over all sources in the
neighborhood of the n -th box

$v_j = v_j + \Phi(\mathbf{y}_j, \mathbf{x}_i)$;

End;

End;

End;

Implementation can be different!
All we need is to get v_j .

Asymptotic Complexity of SLFMM

Assume that:

- By some magic we can easily find neighbors, and lists of points in each box.
- Translation is performed by straightforward $P \times P$ matrix-vector multiplication, where $P(p)$ is the total length of the translation vector. So the complexity of a single translation is $O(P^2)$.
- The source and evaluation points are distributed uniformly, and there are K boxes, with s source points in each box ($s=N/K$). We call s the *grouping* (or *clustering*) parameter.
- The number of neighbors for each box is $O(1)$.

Then Complexity is:

- For Step 1: $O(PN)$
- For Step 2: $O(P^2K^2)$
- For Step 3: $O(PM+Ms)$
- Total: $O(PN+P^2K^2+PM+Ms) = O(PN+P^2K^2+PM+MN/K)$

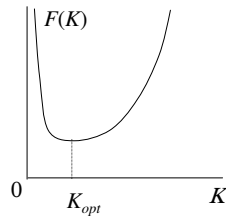
Selection of Optimal K (or s)

$$F(K) = PN + P^2K^2 + PM + PMN/K.$$

$$F'(K) = 2P^2K - PMN/K^2 = 0.$$

$$K_{opt} = \left(\frac{MN}{2P}\right)^{1/3} = O\left(\left(\frac{MN}{P}\right)^{1/3}\right).$$

$$s_{opt} = \frac{N}{K_{opt}} = \left(\frac{2PN^2}{M}\right)^{1/3} = O\left(\frac{PN^2}{M}\right)^{1/3}.$$



Complexity of Optimized SLFMM

$$F(K_{opt}) = PN + P^2\left(\frac{MN}{2P}\right)^{2/3} + PM + PMN\left(\frac{MN}{2P}\right)^{-1/3} \\ = P(M+N) + (MN)^{2/3}O(P^{4/3}).$$

At $K = K_{opt}$, and $M = O(N)$, the complexity of SLFMM is:

$$O(PN + P^{4/3}N^{4/3}) = O(P^{4/3}N^{4/3}).$$

Example of Complexity:

$$P = 10, N = 10^5$$

Straightforward $O(N^2)$: Complexity $\sim 10^{10}$

SLFMM $O((PN)^{4/3})$: Complexity $\sim 10^8$

100 Times CPU savings !

Sorry, but my PC
cannot solve such
a problem!

$$P = 10, N = 10^8$$

Straightforward $O(N^2)$: Complexity $\sim 10^{16}$

SLFMM $O((PN)^{4/3})$: Complexity $\sim 10^{12}$

10000 Times CPU savings !

