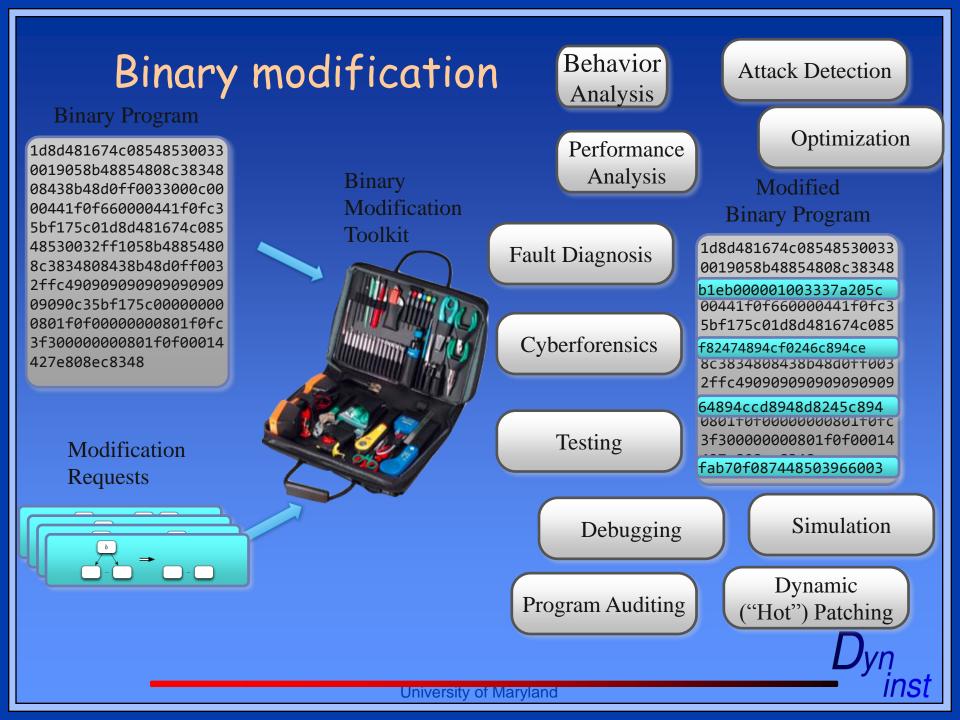
Dyninst: A Binary Analysis and Modification Framework

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Uses For Runtime Code Patching

- Security & Testing
 - Code coverage testing
 - Monitoring (dynamic taint analysis)
- Correctness debugging
 - Fast conditional breakpoints
 - Data breakpoints
- Execution driven simulation
 - Architecture studies



Why Binary Analysis and Manipulation?

- It's what runs on the computer
- All compiled languages (more or less) look the same as a binary
- No Source Code Required
 - For commercial and malware, often not available
- Implicitly Picks up compiler issues
 - Security problems due to compiler bugs



What is Dyninst?

• API for

- binary analysis
- binary re-writing
- runtime patching

• Features

- Generates info about the binary
 - Example: Recover control flow graphs
- New code can be added to programs during execution
 - Permits instrumentation and modification
- Provides processor independent abstractions
- Platform independent patching
 - API abstracts away OS hardware differences

Dyninst Design Philosophy

- Use Any Data Available
 - Debug symbols
 - Dynamic Linker info
 - Binary Analysis within Dyninst
 - User Supplied Info

Work when any source of data is missing

- Stripped binaries
- Static linked program
- Obfuscated binaries



Type & Variable Support in Dyninst

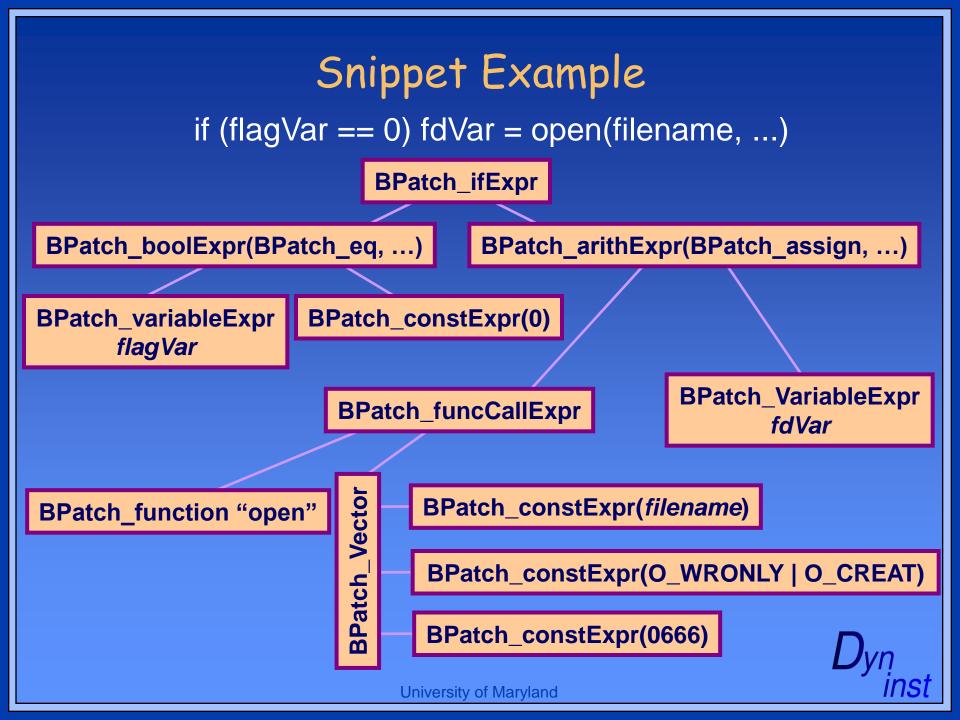
- Access to local (stack) variables
- Complex types
 - non-integer scalars
 - structures
 - arrays
 - Fortran common blocks
- Example: Correctness debugging
 print contents of data structures



Representing New Code Snippets

• Platform Independent Representation

- Same code can be inserted into apps on any system
- Simple Abstract Syntax Tree
 - Can refer to application state (variables & params)
 - Includes simple looping construct
 - Permits calls to application subroutines
- Type Checking
 - Ensures that snippets are type compatible
 - Based on structural equivalence
 - allows flexibility when adding new code



Memory Instrumentation

Dynamic memory access instrumentation

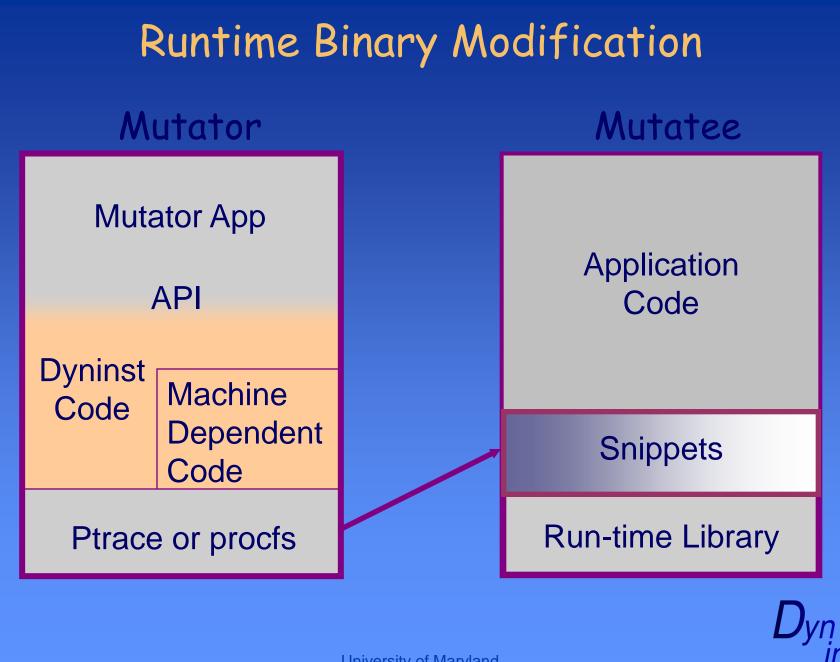
- collect low level memory accesses
- with the flexibility of dynamic instrumentation
- Possible applications
 - tools to catch memory errors
 - offline performance analysis (Sigma etc.)
 - online optimization

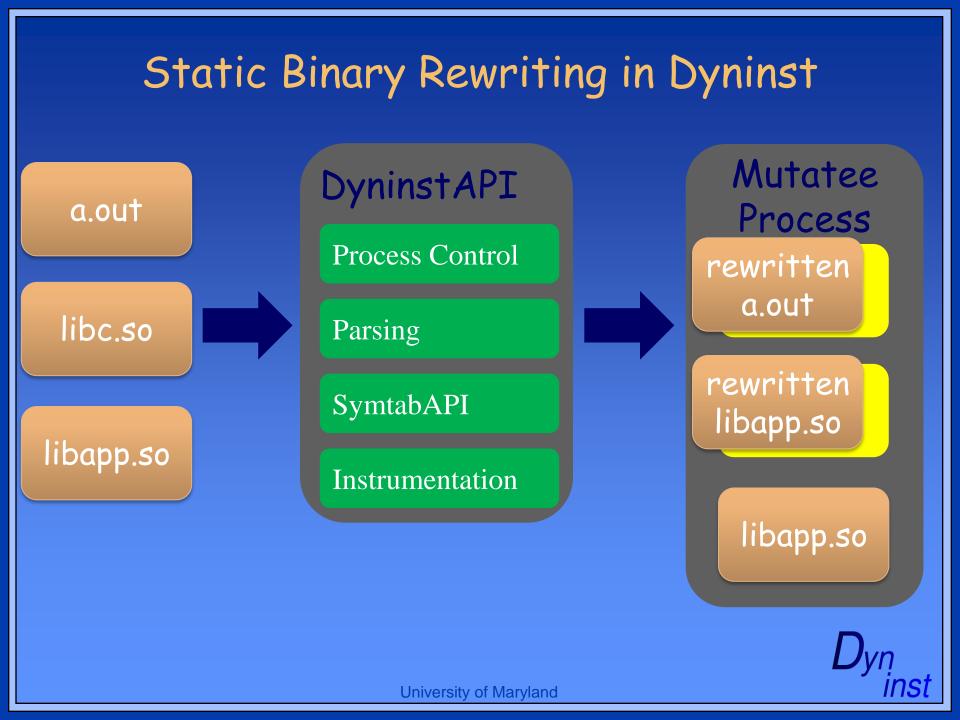


Memory Instrumentation Features

- Finding memory access instructions
 - loads, stores, prefetches
- Builds on Arbitrary Instrumentation
- Decoded instruction information
 - type of instruction
 - constants and registers involved in computing
 - the effective address
 - the number of bytes moved
 - available in the mutator before execution
- Memory access snippets
 - effective address in process space
 - byte count
 - available in mutatee at execution time

Dyn inst





A Static Binary Rewriter

- Binary Rewriter Capabilities
 - Instrument once, run many times
 - Run instrumented binaries on systems without dynamic instrumentation (e.g. some embedded systems).
 - Perform static analysis without running a binary
- Operates on unmodified binaries.
 - No debug information required
 - No linker relocations required
 - No symbols required

• Same abstractions and interfaces as online rewriter.





Static Vs. Dynamic Rewriting

Static Rewriting Faster instrumentation insertion.

✓ Amortize parsing and instrumentation time across multiple runs.

✓Easier to port.

Dynamic Instrumentation ✓Insert and Remove
instrumentation at run time.

 Execute instrumentation at a particular time (oneTimeCode).

Respond to run time events
 (shared library loads, exec, ...).

Binary Rewriting



```
BPatch_addressSpace
• Use BPatch addressSpace for static
 and dynamic code instrumentation.
 if (use bin edit)
  addr space = bpatch.openFile(...);
 else
  addr space = bpatch.attachProcess(...);
```

addr_space->getImage()->findFunction(...); addr_space->insertSnippet(...); addr space->replaceFunction(...);

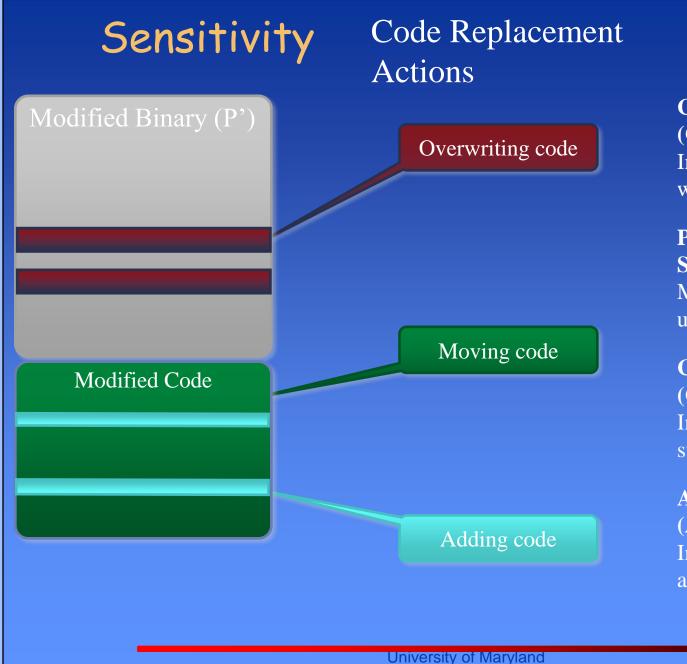


Example Use: Rewriting Symbols Tables Add a function symbol to a binary: /* Open a file */ Symtab *symt; Symtab::openFile(symt, "a.out"); /* Add Symbol */ symt->createFunction("func1" /*name*/, 0x1000 /*offset*/, 100 /*size*/); /* Write new binary */ symt->emit("rewritten.out");

Sensitivity-resistant code relocation

- Preserve visible behavior
- Relationship of input to output
- Identify *sensitive* instructions
- Those whose behavior is changed
- Compensate for *externally sensitive* instructions
- Those whose sensitivity affects visible behavior
- Approach
- Binary analysis (slicing, symbolic execution)
- Code generation
- Runtime checks





Effects

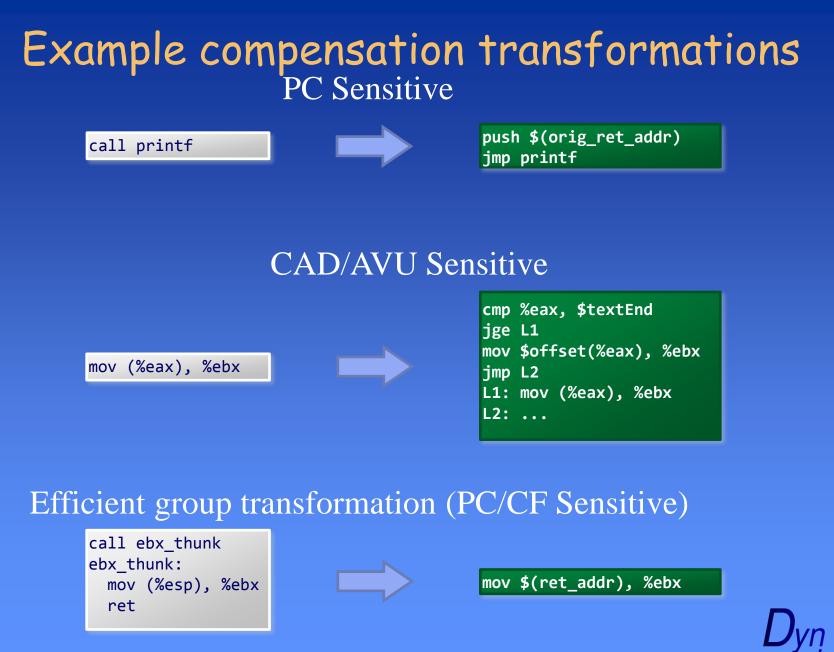
Code-as-Data (CAD) Sensitive Instructions that read or write original code

Program Counter (PC) Sensitive Moved instructions that use the PC

Control Flow (CF) Sensitive Instructions whose successors were moved

Allocated-vs-Unallocated (AVU) Sensitive Instructions that test allocated memory





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Experiments: code relocation

- Verify preservation of behavior on sensitive binaries
- Instrument synthetic malware samples
- Samples should execute with unchanged behavior

• Evaluate overall performance

- Null instrumentation of SPEC CPU 2006 benchmarks, Apache, and MySQL
- Sensitivity-resistant code relocation should reduce overhead
- Group transformations should benefit on Apache/MySQL

Results: behavior preservation

| Packer Tool | Market share | CAD sensitive | Anti-debug | Success |
|------------------|--------------|---------------|------------|---|
| PolyEnE_CAD | 6.21% | yes | | ✓ |
| EXECryptor | 4.06% | yes | yes | |
| Themida | 2.95% | yes | yes | |
| PECompact_CAD | 2.59% | yes | | 1 |
| ASProtect | 0.43% | yes | | 1 |
| Armadillo | 0.37% | yes | yes | |
| Yoda's Protector | 0.33% | yes | yes | Image: A set of the set of the |

S-R relocation succeeded on four additional packers
Failures are due to anti-debug techniques not yet addressed

The Dyninst Team

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Summary

Dyninst Provides

- Multi Architecture Support (x86, Power)
- Multi OS Support (Windows, Linux, AIX, V×Works)
- Multi Compilter (Intel, Microsoft, GCC, PGI, Cray)
- Toolkit approach
 - Uses as little or as much as you want
- Dyninst is Mature
 - Commercial Products from IBM & SGI
 - Used in many third party open source tools
- More Information
 - www.dyninst.org

