

# **Optimum Resource Allocation and Signalling Schemes in Fading CDMA Channels**

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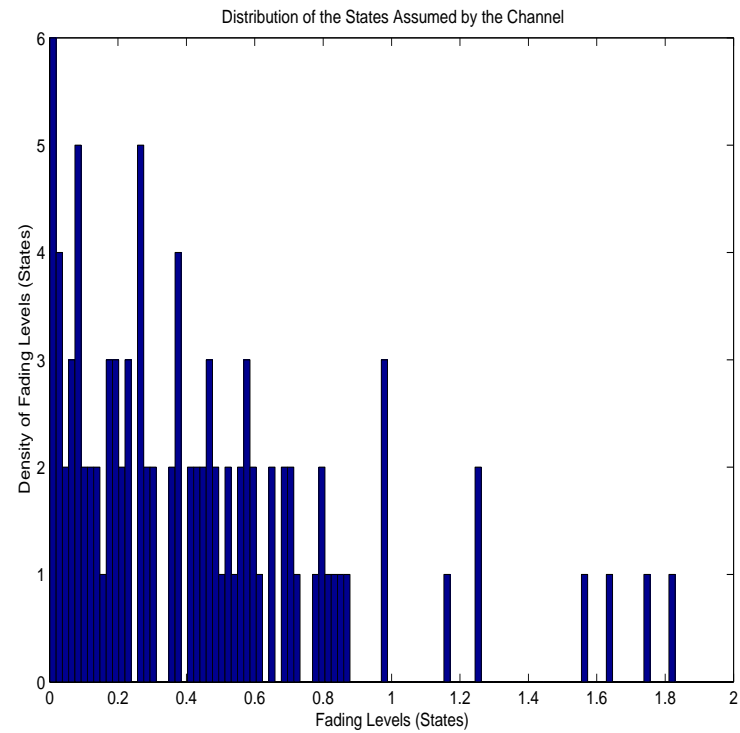
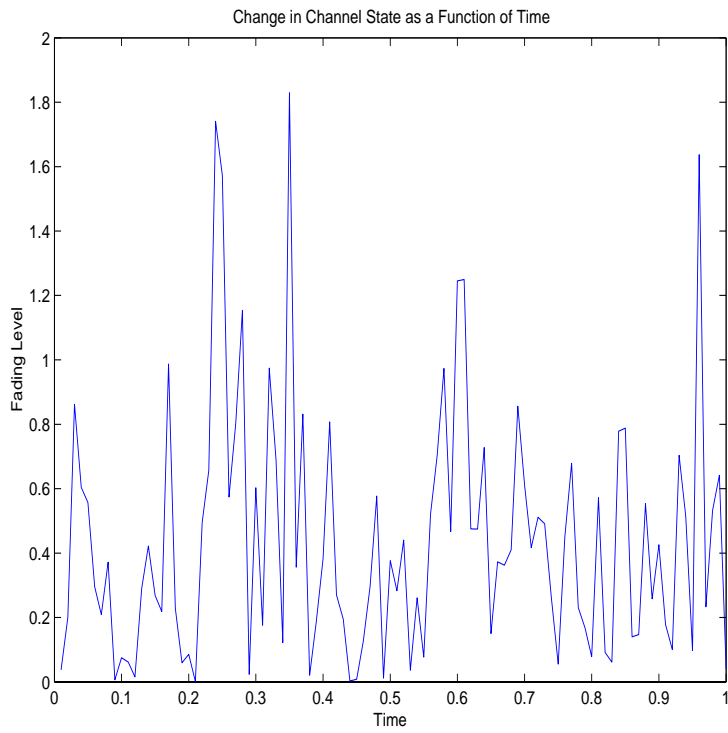
Joint work with Onur Kaya.

## Introduction

- **Fading**: random fluctuations in channel gains.
- If perfect channel state information (CSI) is available at transmitters
  - Dynamic resource allocation to improve quality-of-service or capacity
- Quality-of-service based
  - Provide all users with desired SIR levels
  - Satisfy SIR requirements with minimum transmit power
  - **Compensate** for channel fading; more power if bad channel, less if good channel
- Capacity based
  - Maximize information theoretic ergodic capacity subject to average power constraints
  - **Exploit** variations; more power if good channel, less if bad, no power if very bad

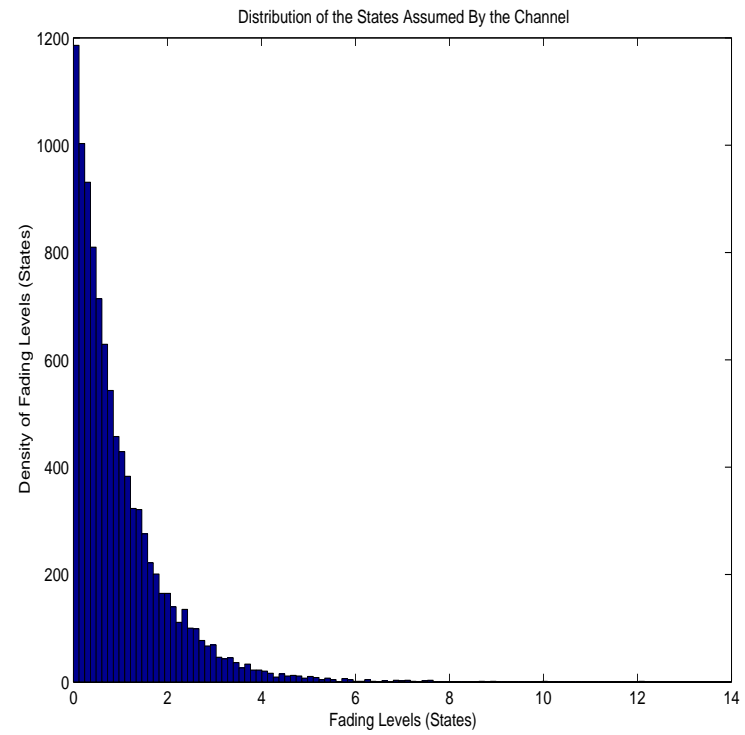
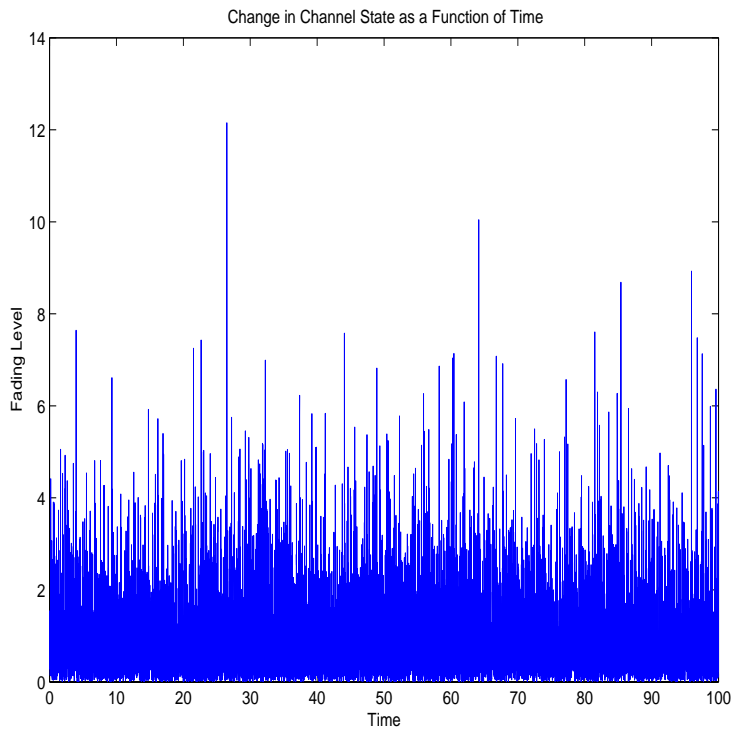
# Illustration of the Channel States

$$r = \sqrt{phx} + n$$



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## Single User Channel (Goldsmith-Varaiya 1994)

- Channel capacity for single user

$$\begin{aligned} C &= \log(1 + SNR) \\ &= \log\left(1 + \frac{P}{\sigma^2}\right) \end{aligned}$$

- In the presence of fading, the capacity for a fixed channel state  $h$ ,

$$C(h) = \log\left(1 + \frac{p(h)h}{\sigma^2}\right)$$

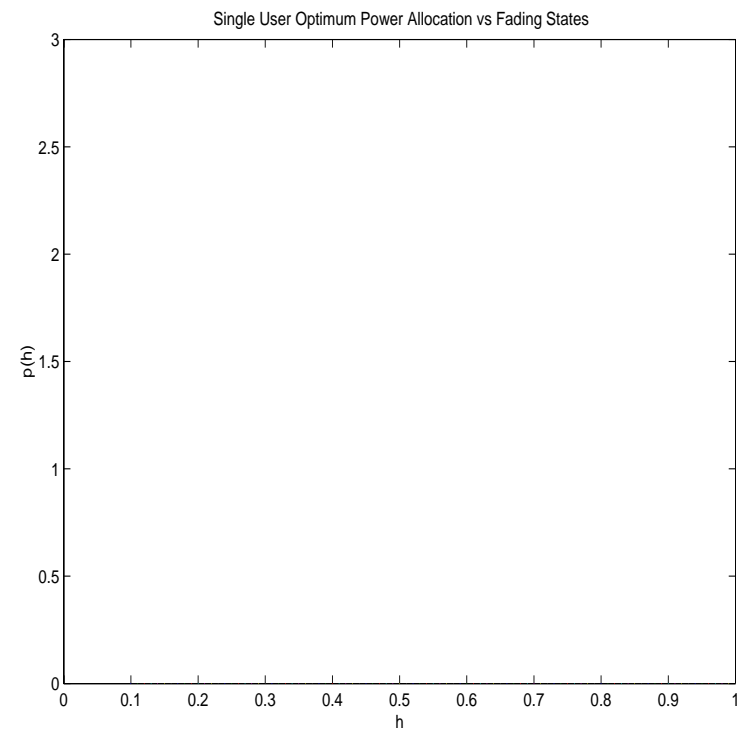
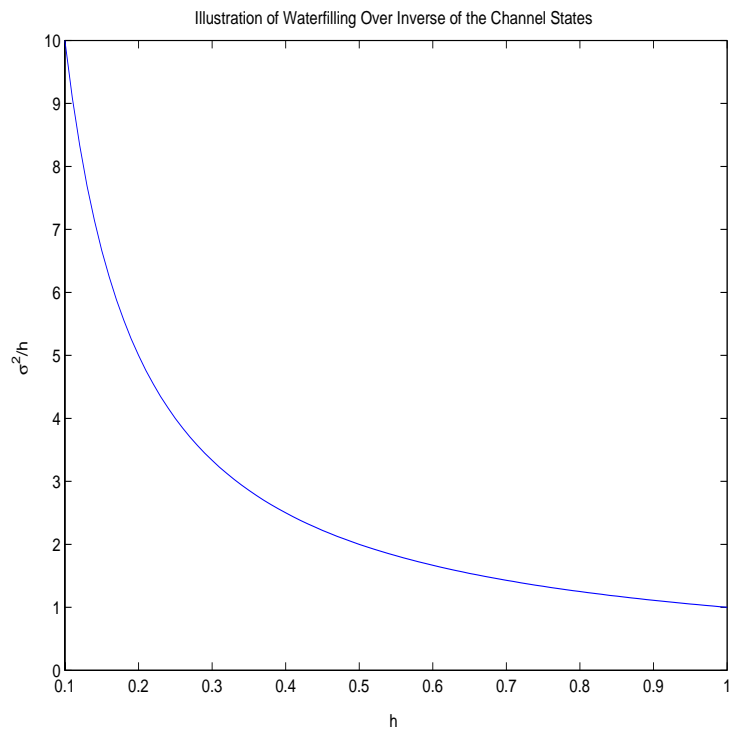
- Maximize the ergodic (expected) capacity, given an average power constraint

$$\begin{aligned} \max_{\{p(h)\}} & E_h \left[ \log\left(1 + \frac{p(h)h}{\sigma^2}\right) \right] \\ \text{s.t.} & E_h [p(h)] \leq \bar{p} \end{aligned}$$

# Single User Channel Solution-Waterfilling

- Optimal power allocation: **waterfilling** of power over time

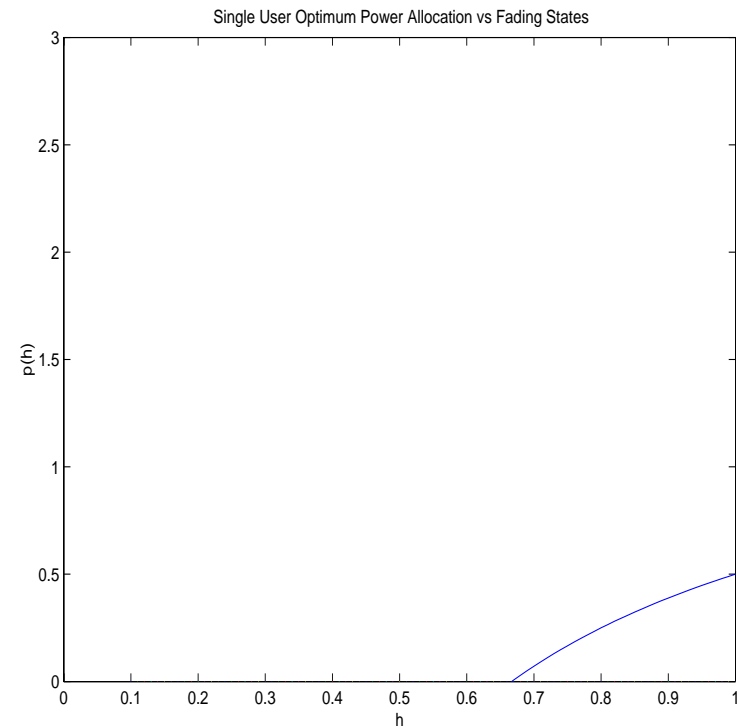
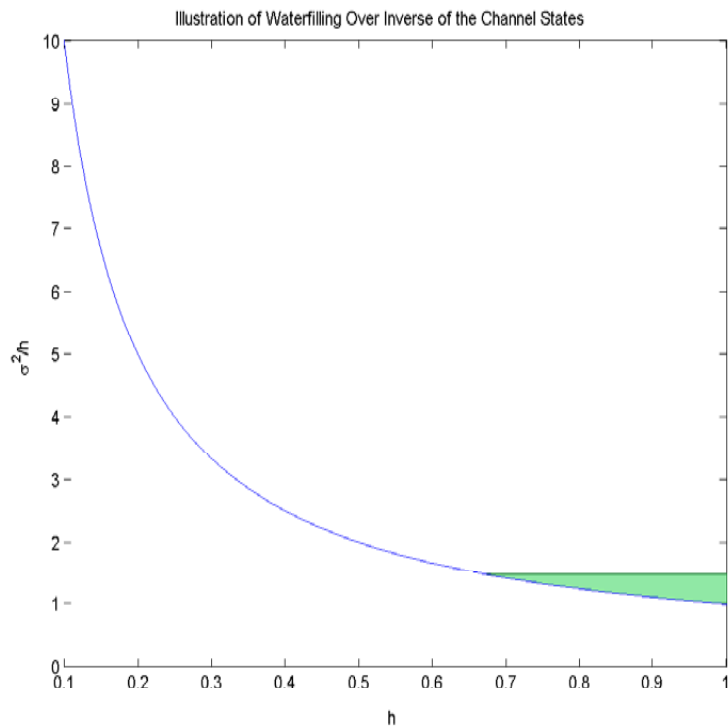
$$p(h) = \left( \frac{1}{\lambda} - \frac{\sigma^2}{h} \right)^+$$



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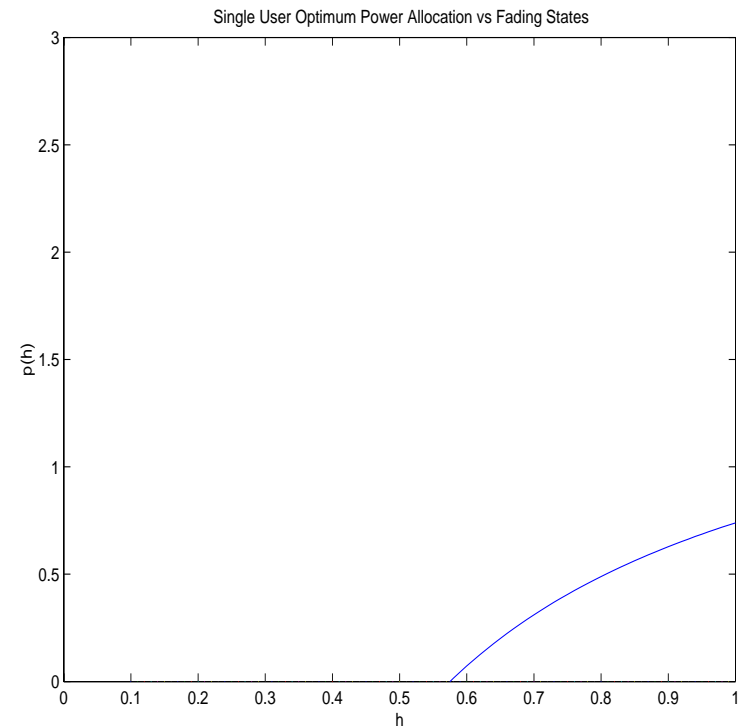
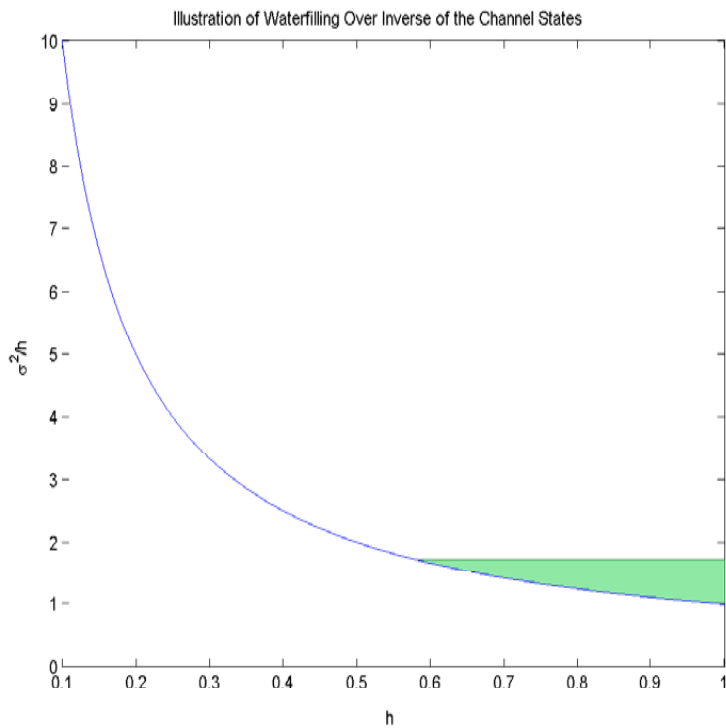
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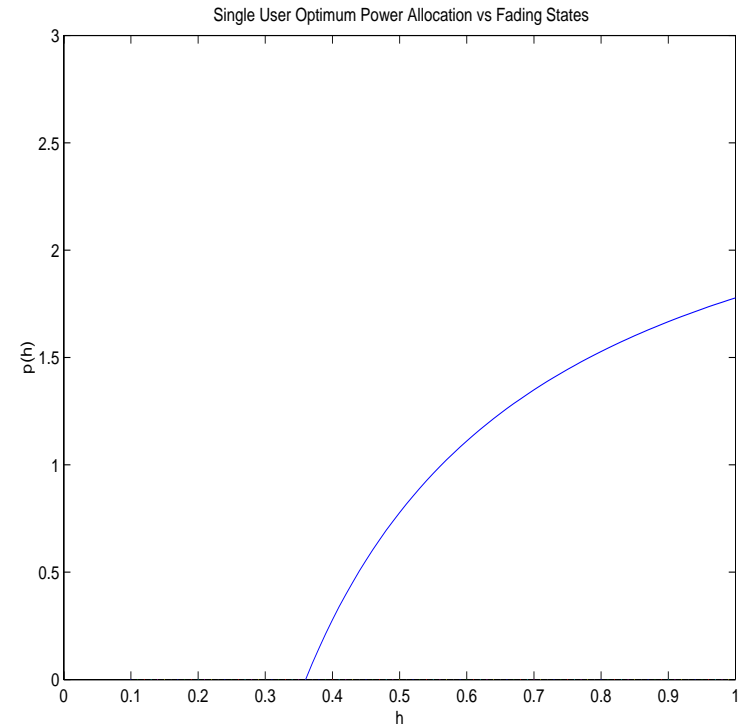
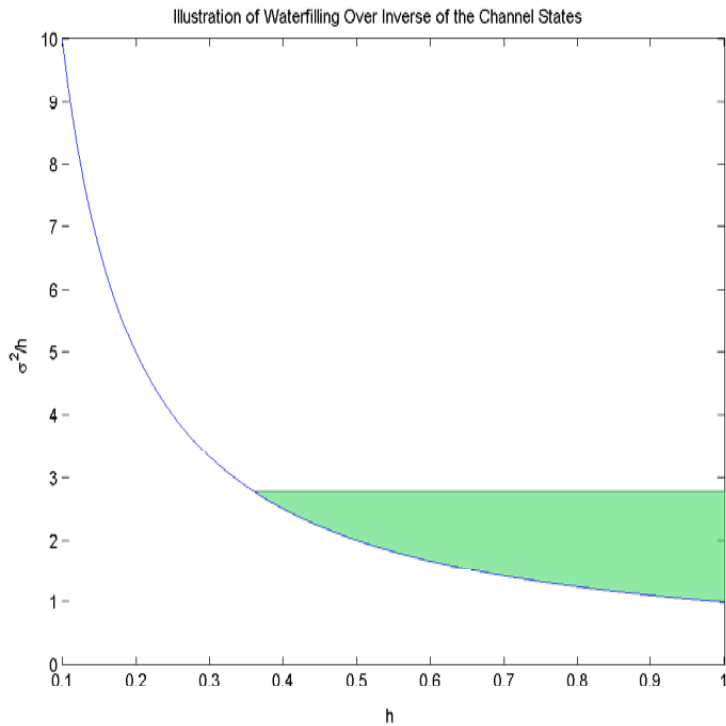




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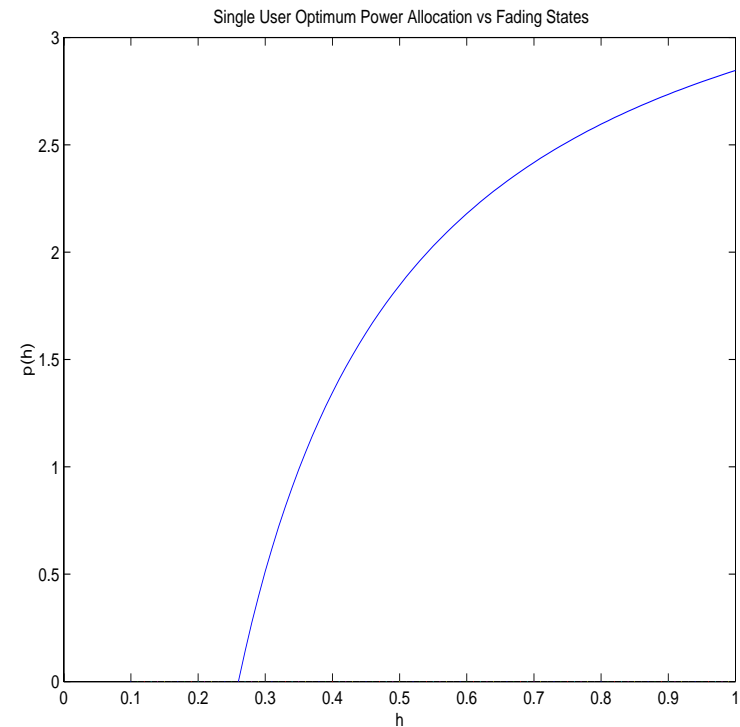
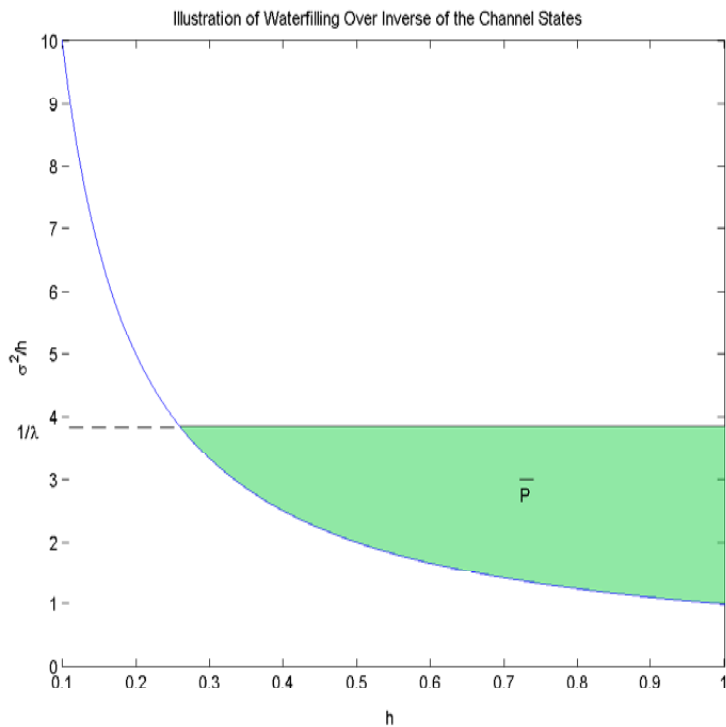
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# Single User Channel Solution-Waterfilling

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$$p(h) = \left( \frac{1}{\lambda} - \frac{\sigma^2}{h} \right)^+$$



## Multiuser Scalar Gaussian Channel (Knopp-Humblet 1995)

- Multiple users, scalar transmissions

$$r = \sum_{i=1}^K \sqrt{p_i(\mathbf{h})} h_i x_i + n$$

- Maximize ergodic **sum capacity**, given average power constraints

$$\begin{aligned} \max_{\{p_i(\mathbf{h})\}} \quad & E_{\mathbf{h}} \left[ \log \left( 1 + \sigma^{-2} \sum_{i=1}^K h_i p_i(\mathbf{h}) \right) \right] \\ \text{s.t.} \quad & E_{\mathbf{h}} [p_i(\mathbf{h})] \leq \bar{p}_i, \quad p_i(\mathbf{h}) \geq 0, \quad i = 1, \dots, K \end{aligned}$$

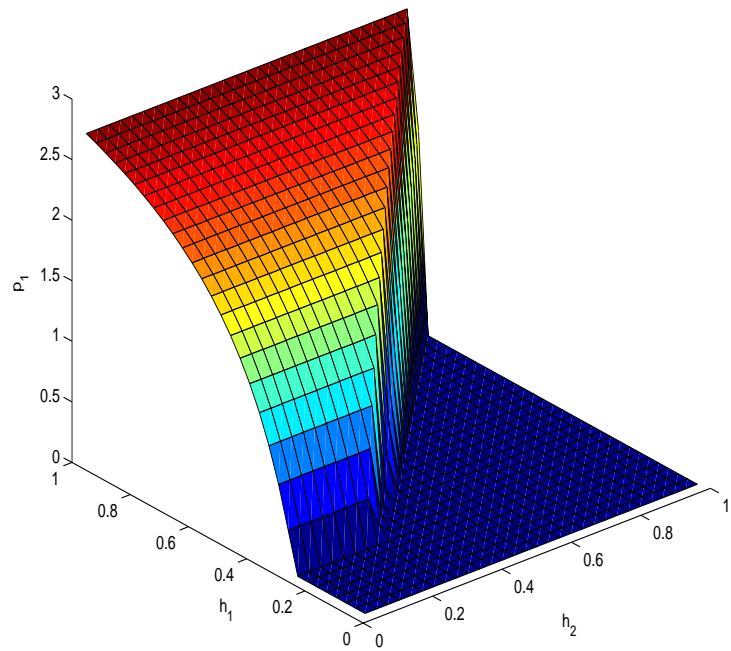
- Optimal power allocation: single user waterfilling on disjoint sets of channel states

$$p_k(\mathbf{h}) = \begin{cases} \left( \frac{1}{\lambda_k} - \frac{\sigma^2}{h_k} \right)^+, & \text{if } h_k/\lambda_k > h_j/\lambda_j, \quad j \neq k \\ 0, & \text{otherwise} \end{cases}$$

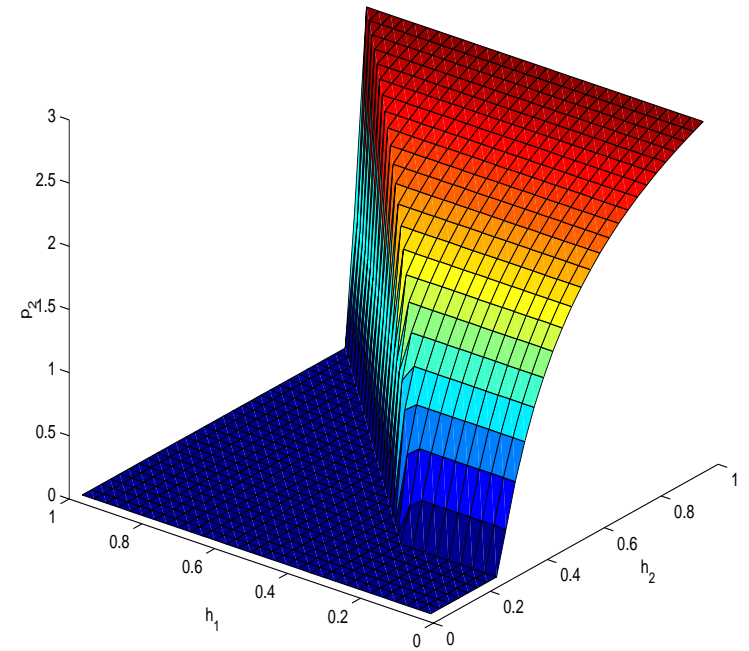
- Only the strongest (after some scaling) user transmits at any given time.

# Optimum Power Allocation: Scalar Multiuser Channel

Power Distribution of User 1



Power Distribution of User 2



## Multiuser Vector (Waveform) Gaussian Channel

- Project the received signal onto  $N$  basis waveforms.
- CDMA: vector signals modulated by scalar symbols.

$$\mathbf{r} = \sum_{i=1}^K \sqrt{p_i(\mathbf{h})} h_i x_i \mathbf{s}_i + \mathbf{n}$$

- Maximize ergodic **sum capacity** subject to average power constraints

$$\begin{aligned} \max_{\{\mathbf{p}(\mathbf{h})\}} & E_{\mathbf{h}} \left[ \log \left| \mathbf{I}_N + \sigma^{-2} \sum_{i=1}^K h_i p_i(\mathbf{h}) \mathbf{s}_i \mathbf{s}_i^{\top} \right| \right] \\ \text{s.t.} & E_{\mathbf{h}}[p_i(\mathbf{h})] \leq \bar{p}_i, \quad i = 1, \dots, K \\ & p_i(\mathbf{h}) \geq 0, \quad \forall \mathbf{h}, \quad i = 1, \dots, K \end{aligned}$$

## Optimal Power Control

- $C_{\text{sum}}$  is a concave function of powers. Constraint set is convex.
- Using Lagrange method, optimum powers satisfy (by KKT conditions),

$$\frac{h_k \mathbf{s}_k \mathbf{A}_k^{-1} \mathbf{s}_k}{1 + p_k(\mathbf{h}) h_k \mathbf{s}_k \mathbf{A}_k^{-1} \mathbf{s}_k} \leq \lambda_k, \quad k = 1, \dots, K, \quad \forall \mathbf{h} \in R^K$$

with equality iff  $p_k > 0$ . Here,  $\mathbf{A}_k$  is defined as

$$\mathbf{A}_k = \sigma^2 \mathbf{I}_N + \sum_{i \neq k} h_i p_i(\mathbf{h}) \mathbf{s}_i \mathbf{s}_i^\top$$

- Optimum power allocation:

$$p_k(\mathbf{h}) = \left( \frac{1}{\lambda_k} - \frac{1}{h_k \mathbf{s}_k^\top \mathbf{A}_k^{-1} \mathbf{s}_k} \right)^+, \quad k = 1, \dots, K$$

- **Simultaneous waterfilling** of powers onto  
inverse of the “SIRs with MMSE receivers and unit transmit powers” of users.

## Iterative Waterfilling

- Isolate  $k$ th user's contribution to sum capacity

$$C_{\text{sum}} = C_k + \bar{C}_k$$

$$C_k = E_{\mathbf{h}} \left[ \log \left( 1 + h_k p_k(\mathbf{h}) \mathbf{s}_k^\top \mathbf{A}_k^{-1} \mathbf{s}_k \right) \right]$$

- Optimize the power of user  $k$  only, with the powers of all other users fixed.

$$\begin{aligned} p_k^{n+1} &= \arg \max_{p_k} C_{\text{sum}}(p_1^{n+1}, \dots, p_{k-1}^{n+1}, p_k, p_{k+1}^n, \dots, p_K^n) \\ &= \arg \max_{p_k} C_k(p_k) \end{aligned}$$

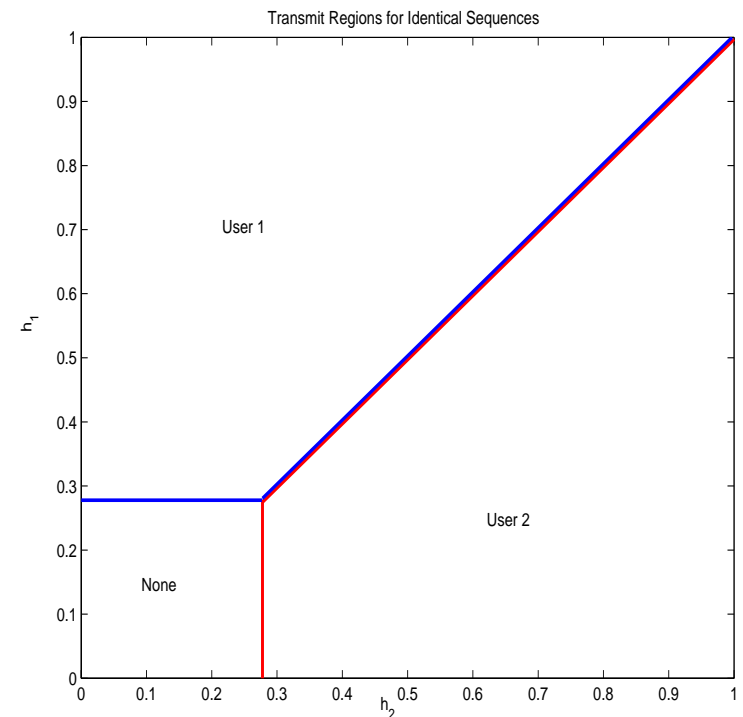
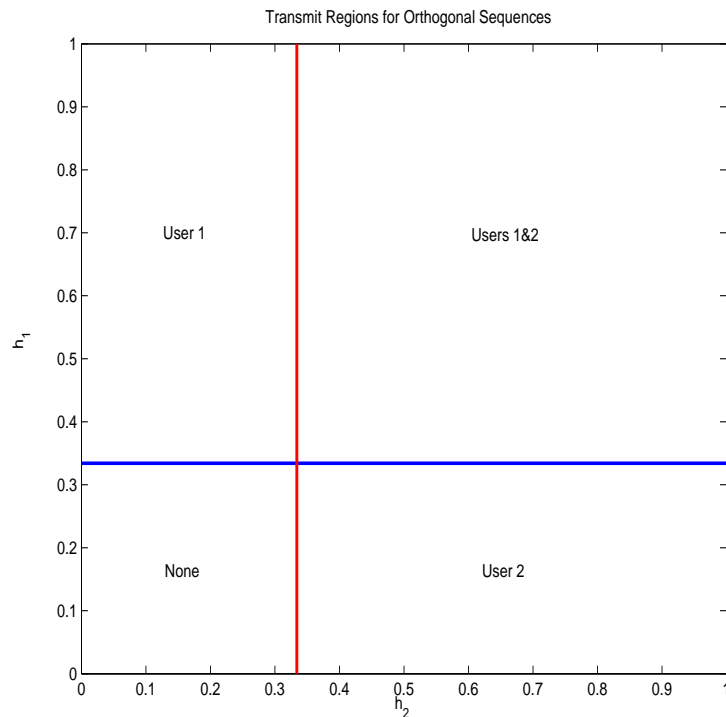
- **One-user-at-a-time** single user waterfilling:

$$p_k(\mathbf{h}) = \left( \frac{1}{\tilde{\lambda}_k} - \frac{1}{h_k \mathbf{s}_k^\top \mathbf{A}_k^{-1} \mathbf{s}_k} \right)^+$$

- Converges to global optimum [Bertsekas-Tsitsiklis].

# Simultaneous Transmit Regions

- The regions where both users transmit for the two special cases:



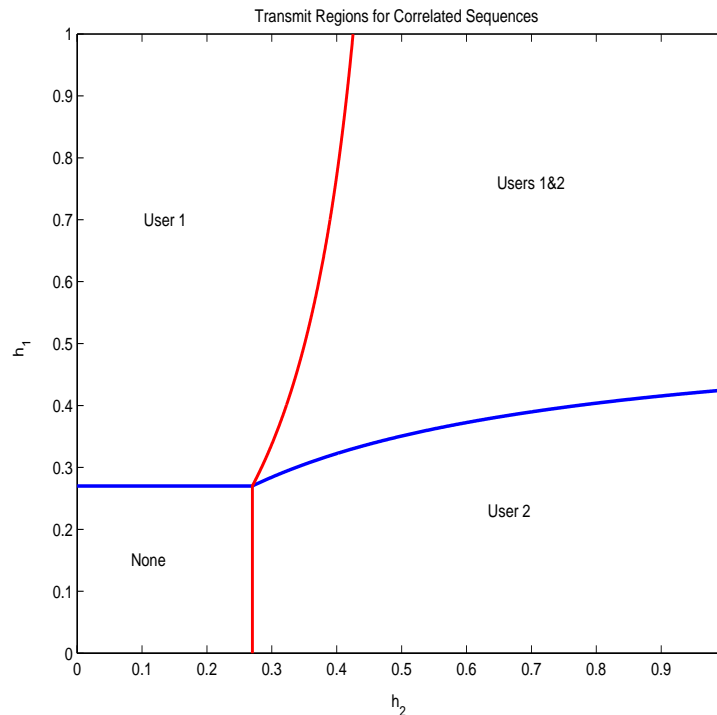
- **Motivation:** for a set of arbitrary signature sequences, is there a set of channel states (with non-zero probability measure) where all users transmit simultaneously?



## Simultaneous Transmit Condition

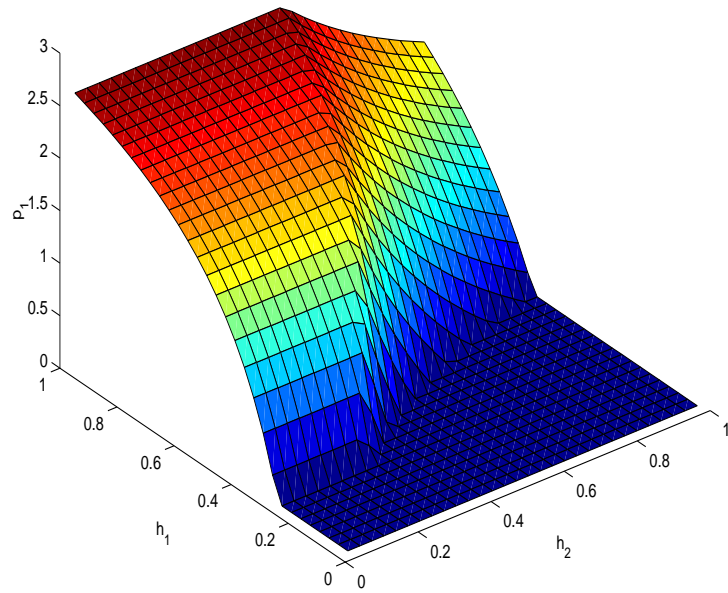
**Theorem:** There exists a non-zero probability region of fading states  $\mathbf{h}$  where all  $K$  users in the system transmit simultaneously, if and only if  $\{\mathbf{s}_i \mathbf{s}_i^\top\}_{i=1}^K$  are linearly independent.

**Corollary:** When  $K \leq N$ , for a set of  $K$  linearly independent signature sequences, there always exists a non-zero probability region of channel states where all  $K$  users transmit simultaneously.

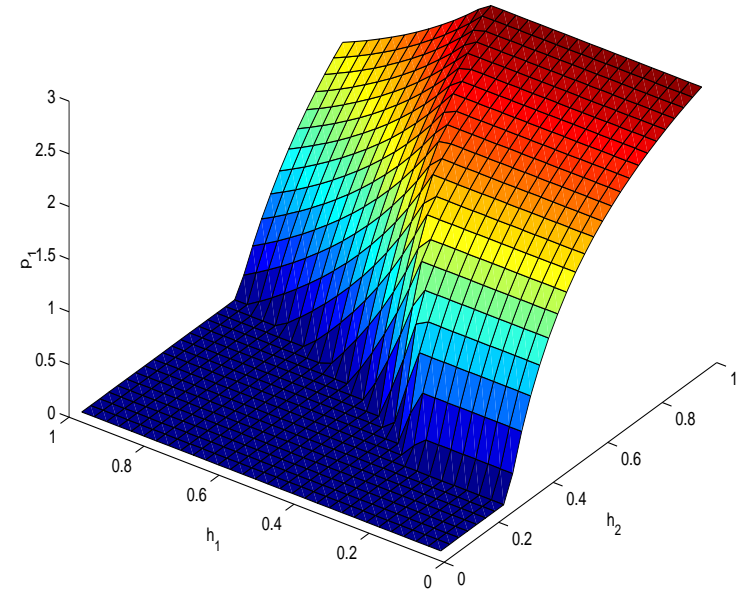


# Transmit Powers: Correlated Signatures

Power Distribution of User 1



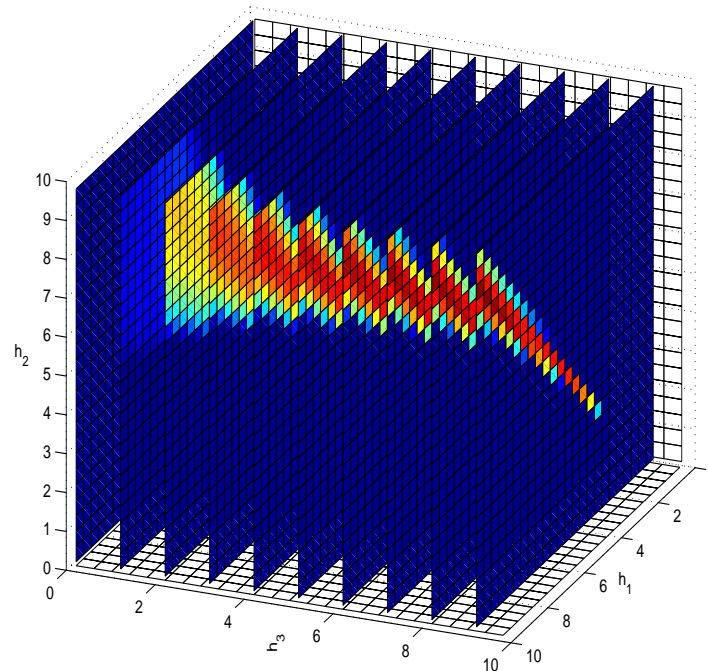
Power Distribution of User 2



## Maximum Number of Simultaneous Transmissions

**Corollary:** For a set of signature sequences with  $\text{rank}(\mathbf{S}) = M \leq \min\{N, K\}$ , the number of users that can transmit simultaneously cannot be larger than  $M(M+1)/2$ .

Example:  $N = 2, K = 3$ .



Signature sequences  $\{\mathbf{s}_i\}_{i=1}^K$  are linearly dependent, but  $\{\mathbf{s}_i \mathbf{s}_i^\top\}_{i=1}^K$  are linearly independent.

## Jointly Optimal Power and Waveform Allocation in Fading

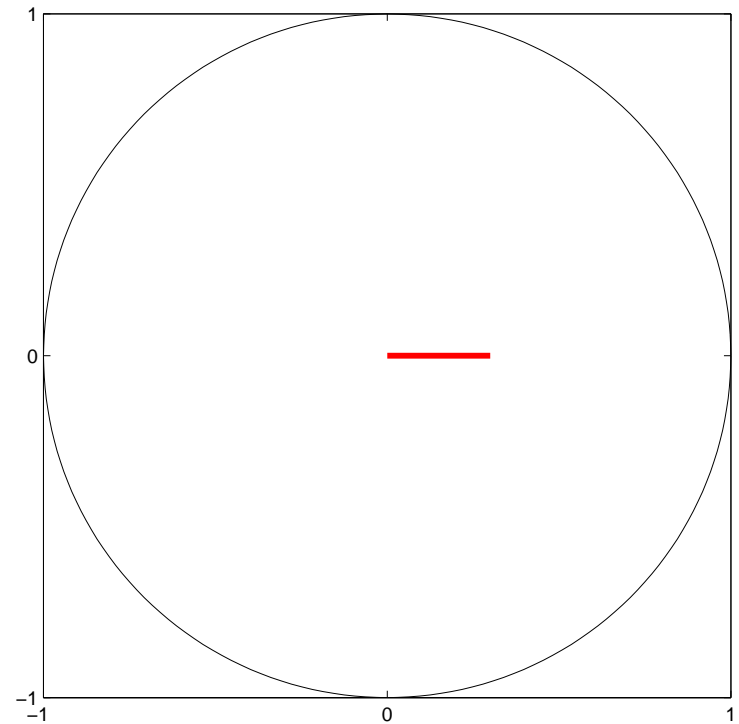
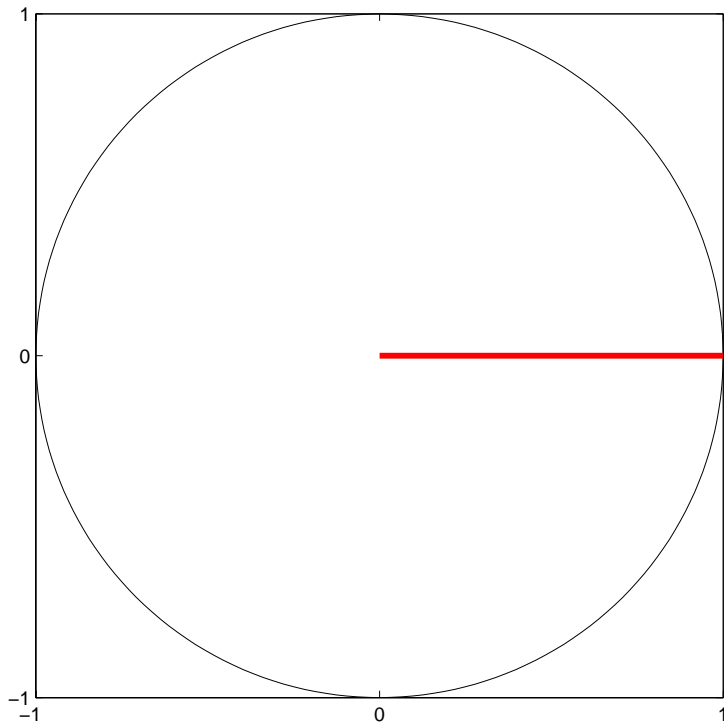
- Dynamic resource allocation – transmit powers, bandwidth, time slots; or in general waveforms – to combat fading and improve capacity
- Vector (waveform) MAC: allocate transmit powers and waveforms to users.

$$\mathbf{r} = \sum_{i=1}^K \sqrt{p_i h_i} x_i \mathbf{s}_i + \mathbf{n}$$

- Existing literature:
  - **Power control only**: control powers as a function of CSI in **fading** [Kaya-Ulukus].
    - \* maximize sum capacity,
    - \* achieve any point on the capacity region (maximize weighted sum of rates).
  - **Waveform allocation only**: find sum-capacity maximizing set of waveforms for a given set of (fixed) powers in **no fading** [Rupf-Massey, Viswanath-Anantharam].
    - \* notion of oversized/non-oversized users according to powers,
    - \* orthogonal waveforms to oversized users, GWBE waveforms to non-oversized users.

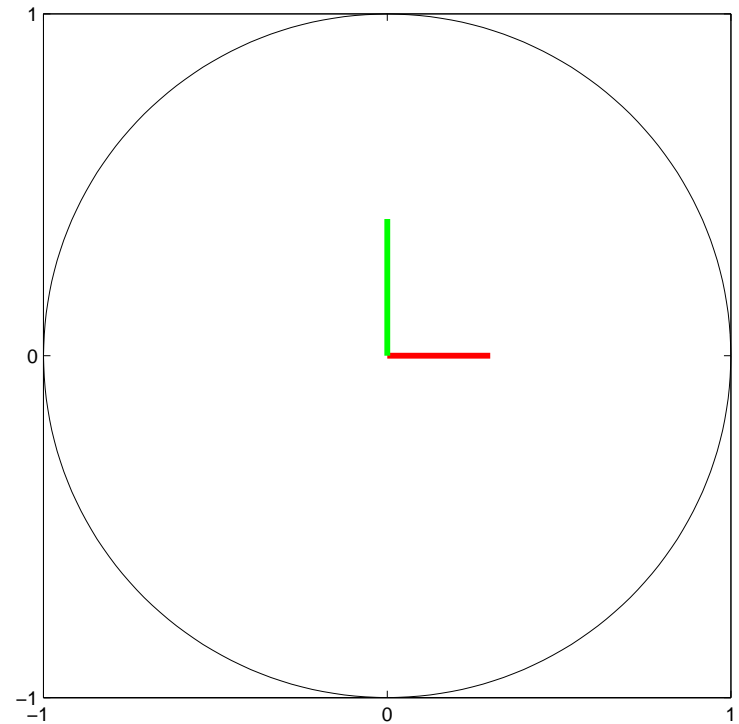
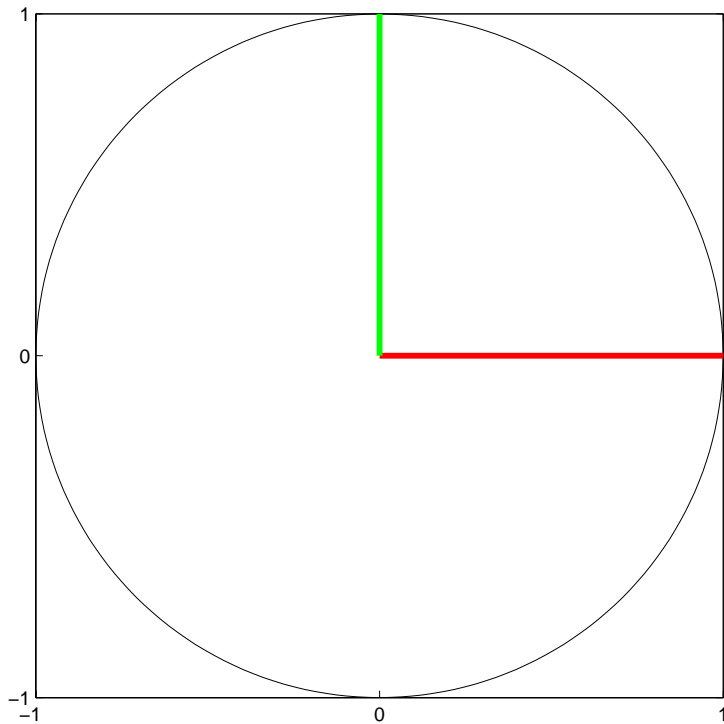
## Waveform Allocation Only – No Fading, Fixed Powers

Simple example: vectors are signatures with powers.



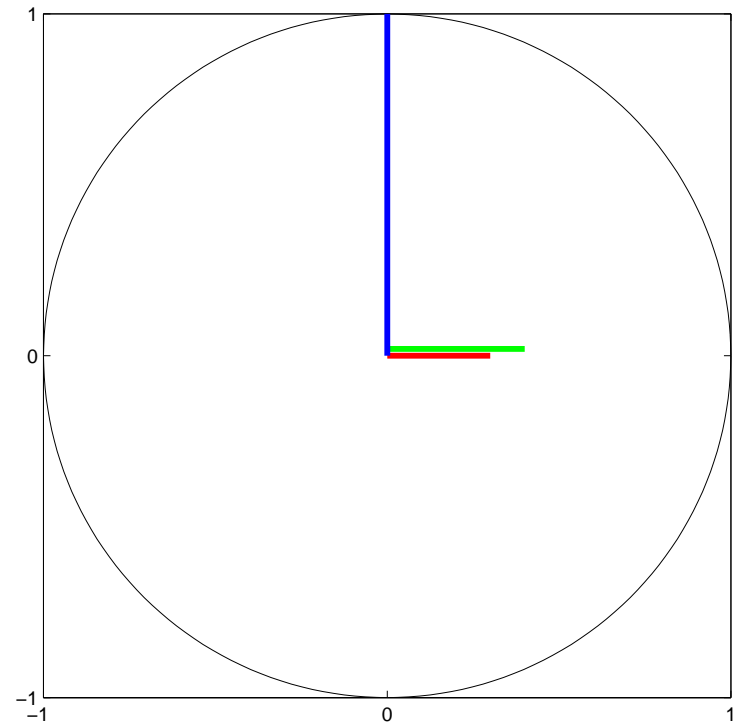
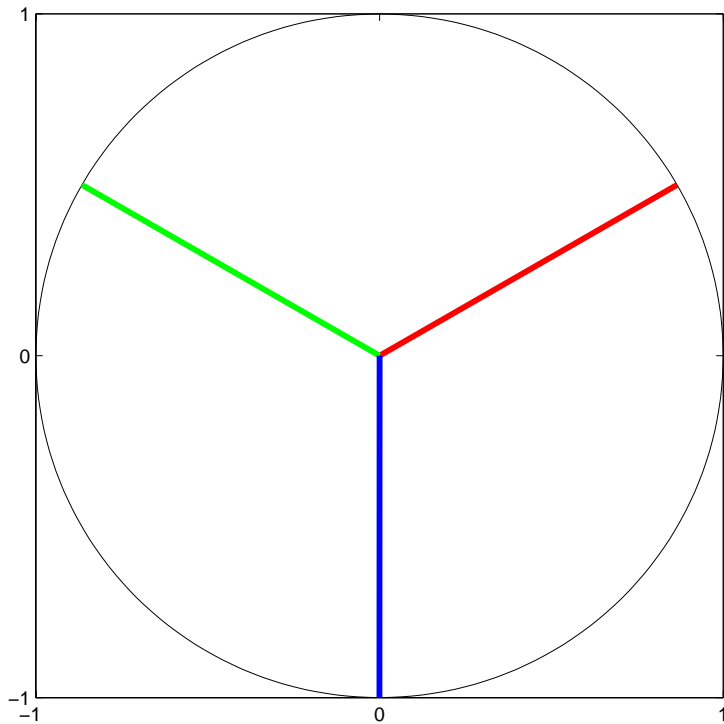
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## Joint Power and Waveform Allocation

- Consider sum capacity of the network. Perfect CSI at the transmitters.
- Then, both powers and waveforms can be chosen as functions of channel states.

$$\mathbf{r} = \sum_{i=1}^K \sqrt{p_i(\mathbf{h})} h_i x_i \mathbf{s}_i(\mathbf{h}) + \mathbf{n}$$

- Ergodic sum capacity maximization problem becomes

$$\begin{aligned} & \max_{\mathbf{p}(\mathbf{h}), \mathbf{S}(\mathbf{h})} && E_{\mathbf{h}} \left[ \log \left| \mathbf{I}_N + \sigma^{-2} \sum_{i=1}^K h_i p_i(\mathbf{h}) \mathbf{s}_i(\mathbf{h}) \mathbf{s}_i(\mathbf{h})^{\top} \right| \right] \\ & \text{s.t.} && E_{\mathbf{h}} [p_i(\mathbf{h})] = \bar{p}_i, \quad i = 1, \dots, K \\ & && p_i(\mathbf{h}) \geq 0, \quad \forall \mathbf{h}, \quad i = 1, \dots, K \\ & && \mathbf{s}_i(\mathbf{h})^{\top} \mathbf{s}_i(\mathbf{h}) = 1, \quad \forall \mathbf{h}, \quad i = 1, \dots, K \end{aligned}$$



## Waveform Optimized Capacity

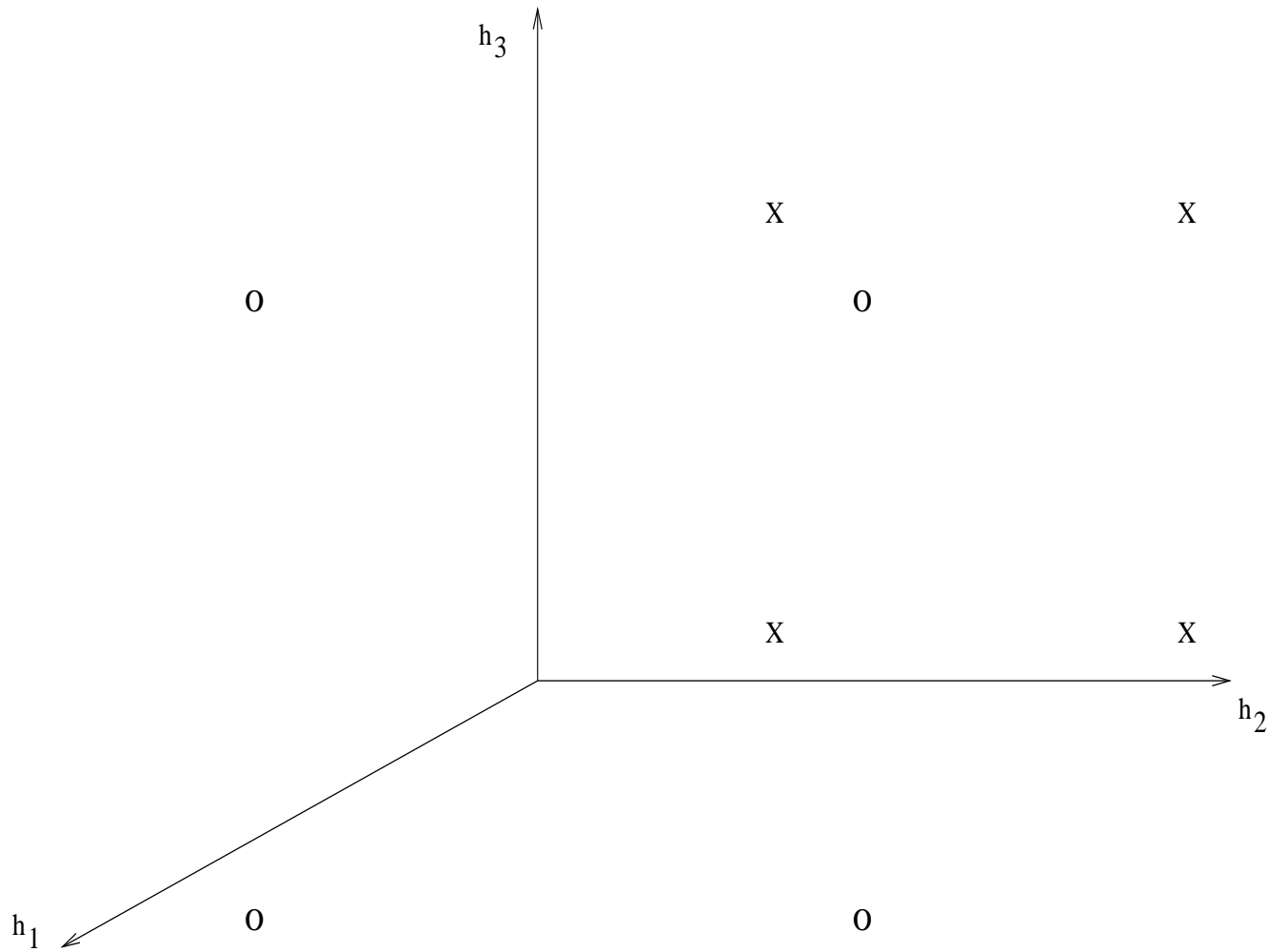
- First, fix an arbitrary valid power allocation over the fading states.
- For each fixed allocation, find the waveforms that maximize the sum capacity at each state  $\mathbf{h}$ .
- Define the waveform-optimized sum capacity at  $\mathbf{h}$

$$C_{\text{opt}}(\mathbf{h}, \mathbf{p}(\mathbf{h})) \triangleq \max_{\mathbf{S}(\mathbf{h})} C_{\text{sum}}(\mathbf{h}, \mathbf{p}(\mathbf{h}), \mathbf{S}(\mathbf{h}))$$

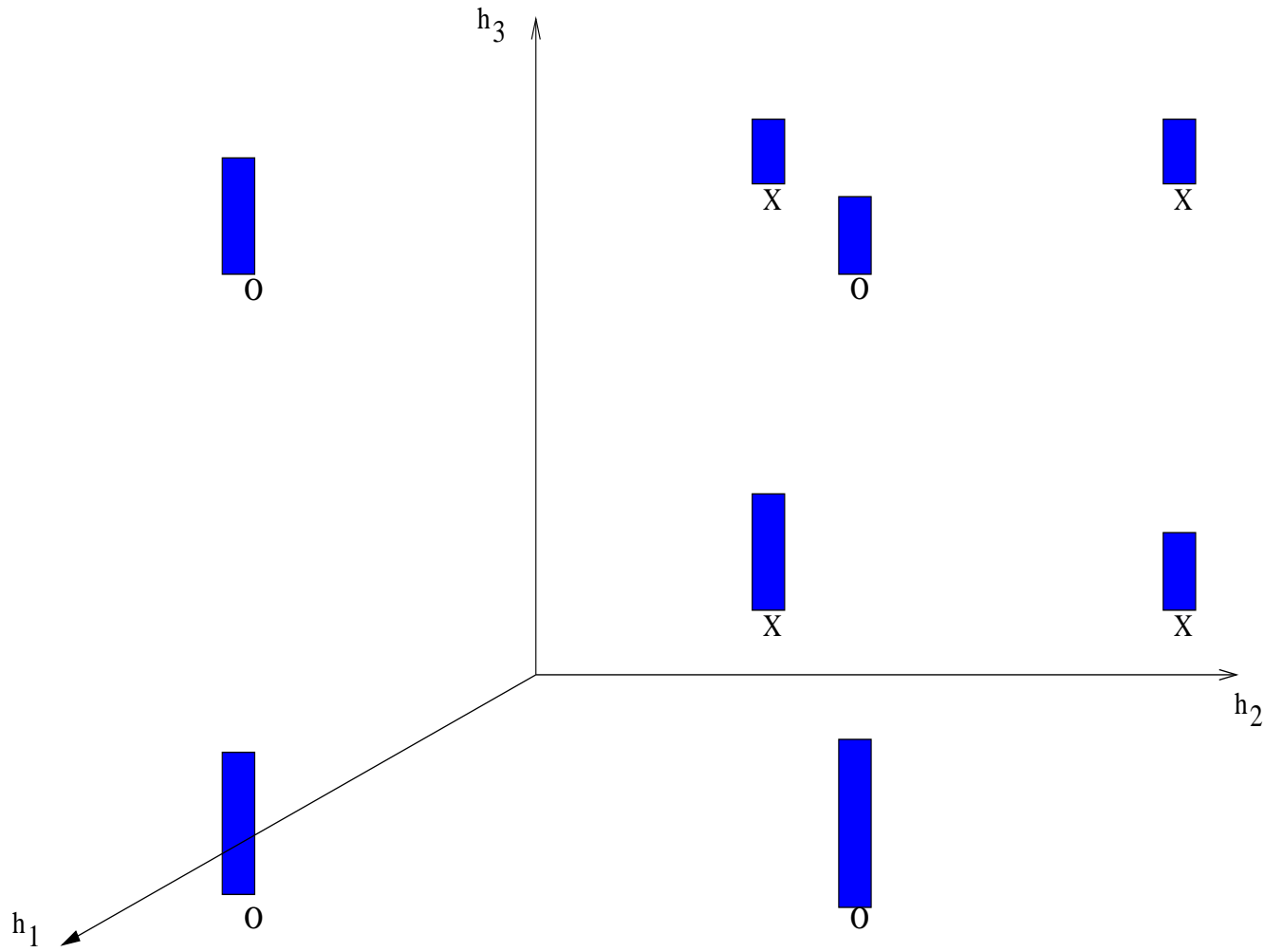
- Then, optimize waveform-optimized sum capacity in terms of the powers,

$$\begin{aligned} \max_{\mathbf{p}(\mathbf{h})} \quad & E_{\mathbf{h}} [C_{\text{opt}}(\mathbf{h}, \mathbf{p}(\mathbf{h}))] \\ \text{s.t.} \quad & E_{\mathbf{h}} [p_i(\mathbf{h})] = \bar{p}_i, \quad i = 1, \dots, K \\ & p_i(\mathbf{h}) \geq 0, \quad \forall \mathbf{h}, \quad i = 1, \dots, K \end{aligned}$$

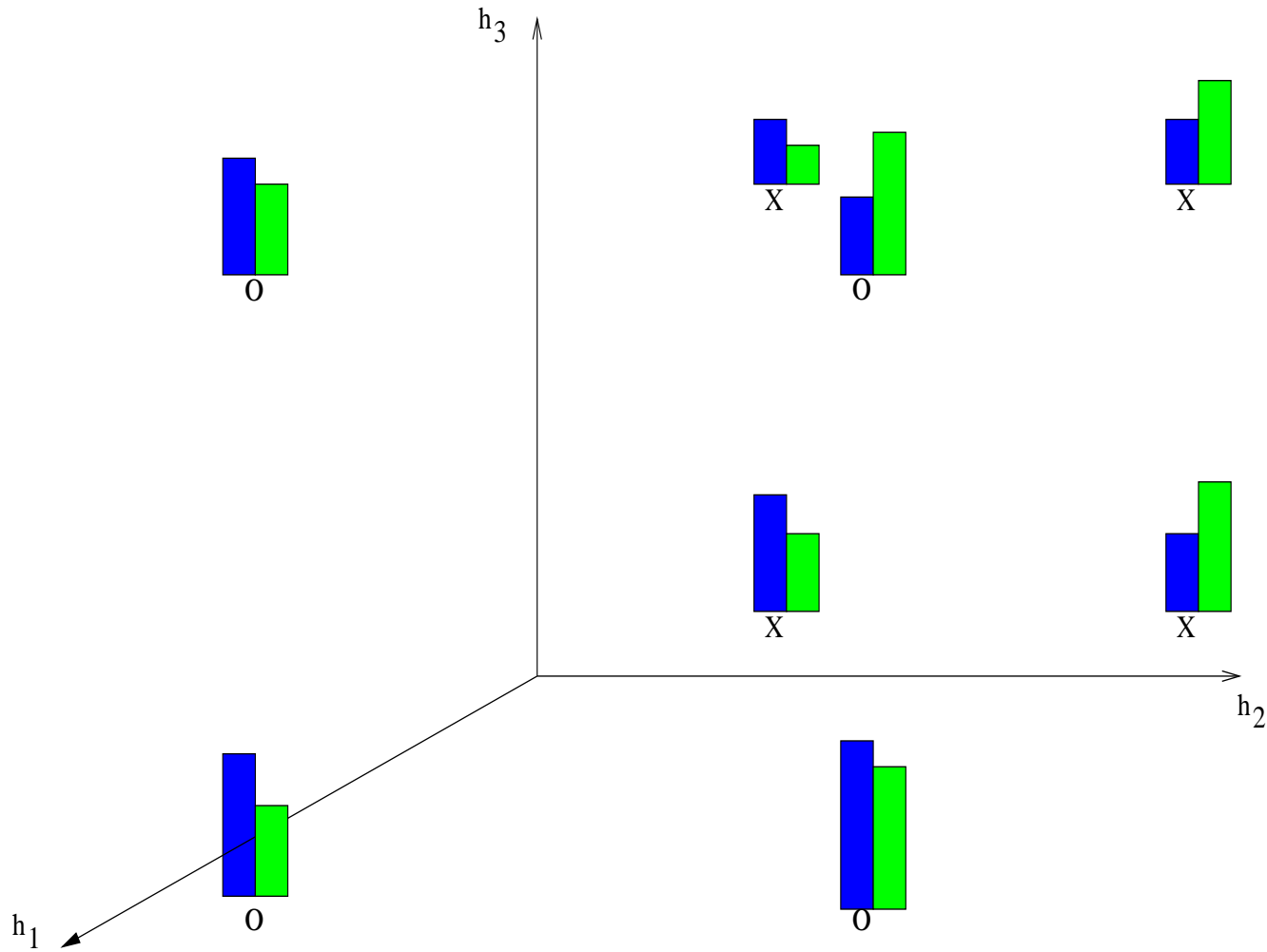
# Choosing the Optimum Waveforms – Illustration



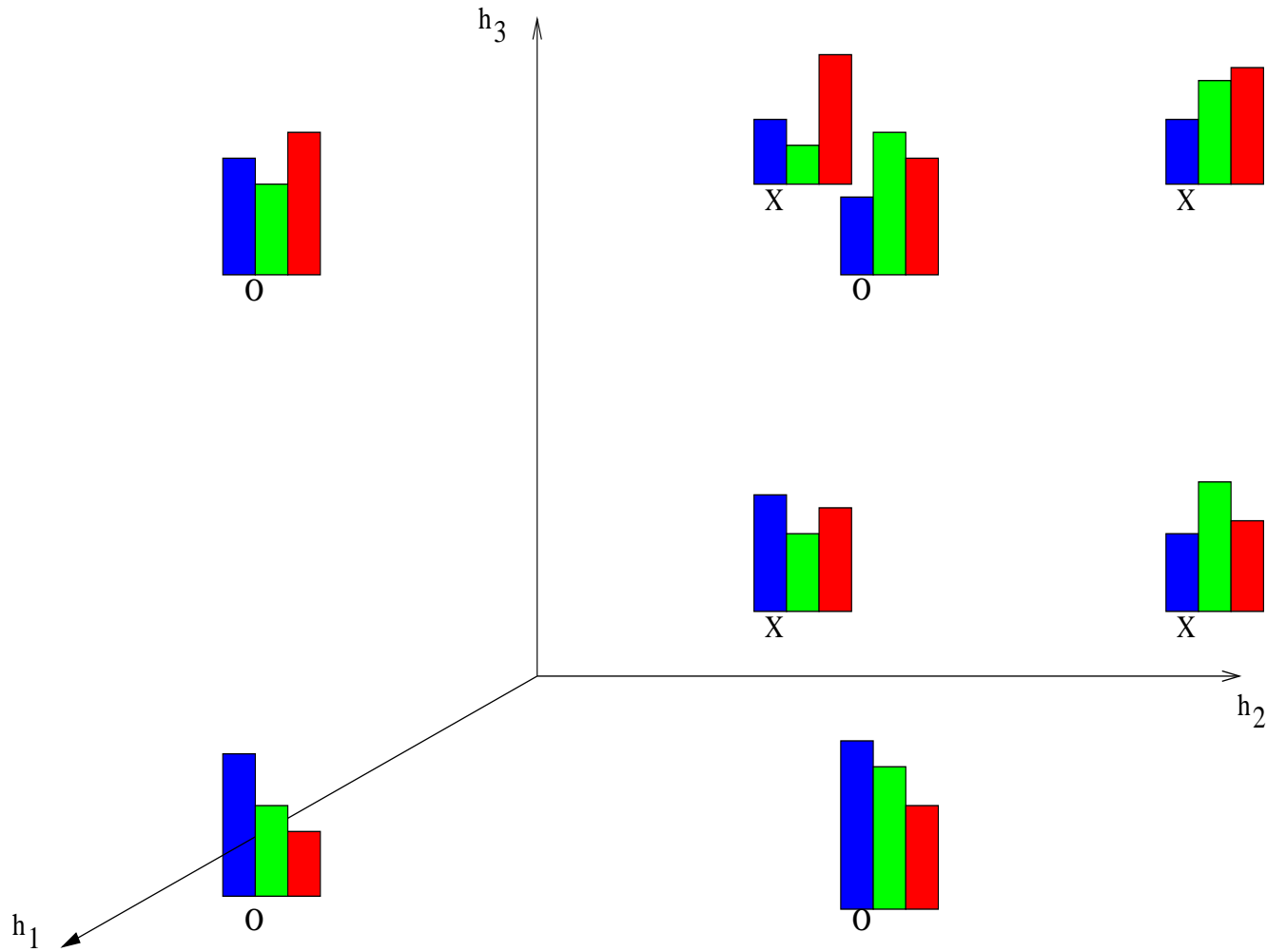
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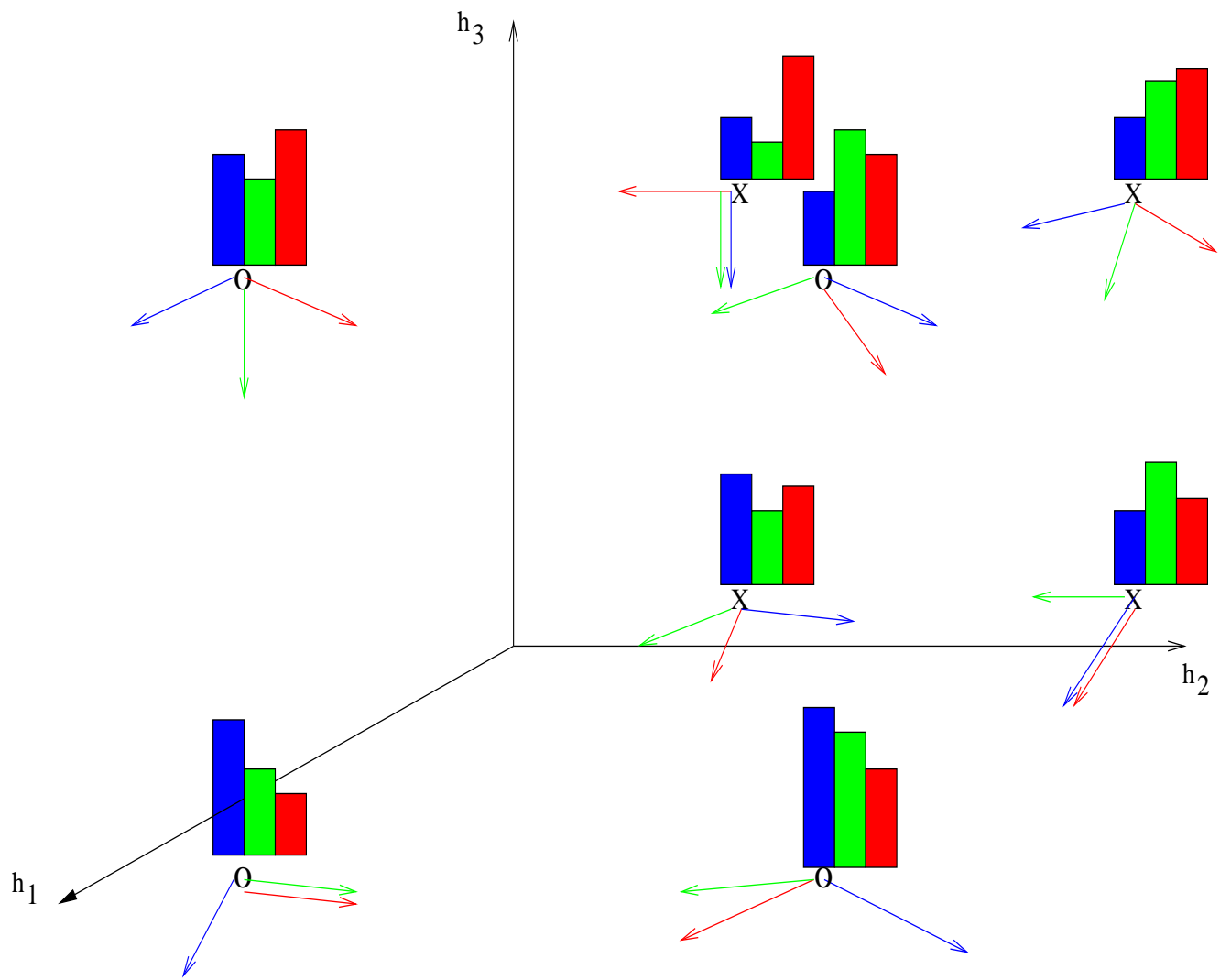
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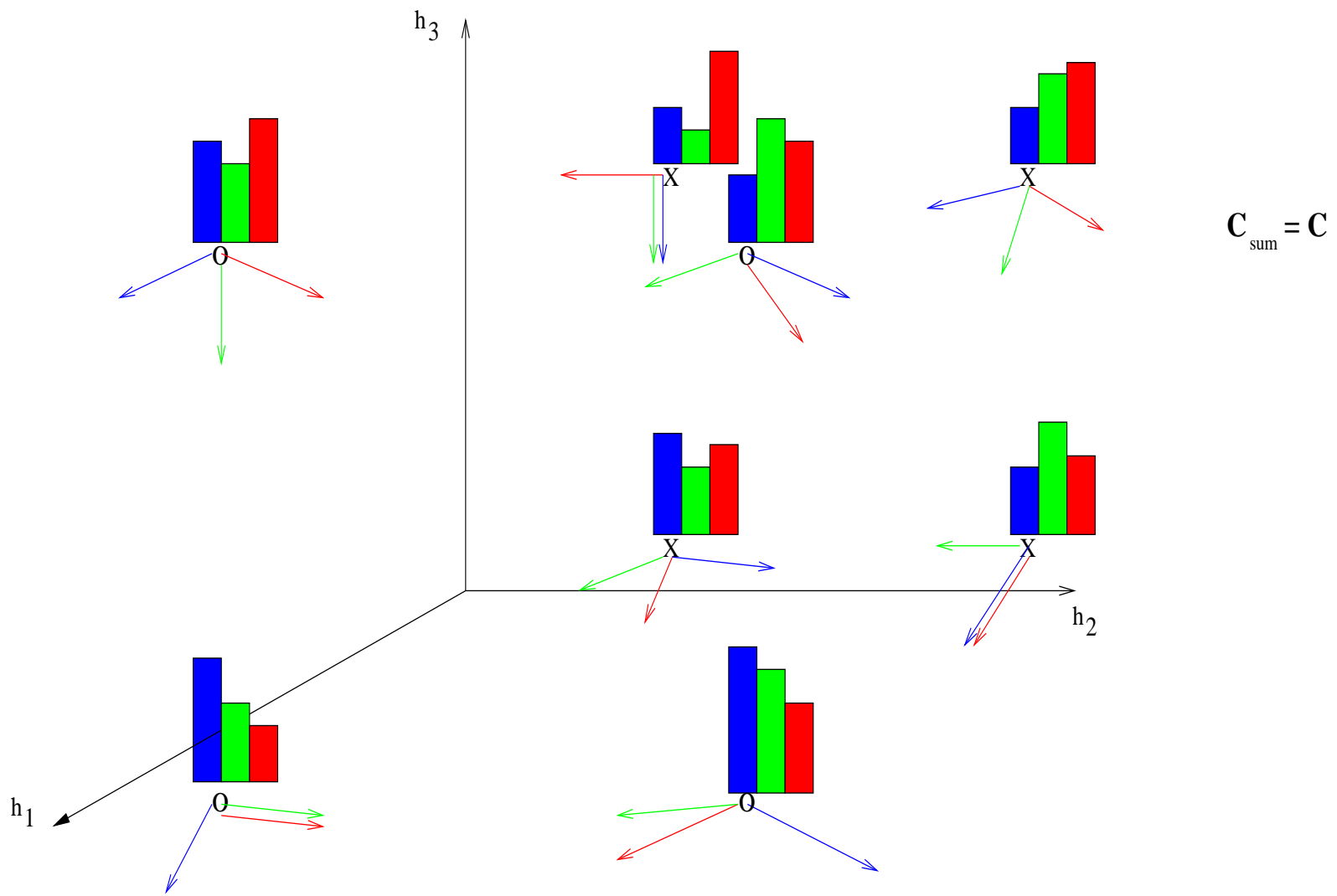
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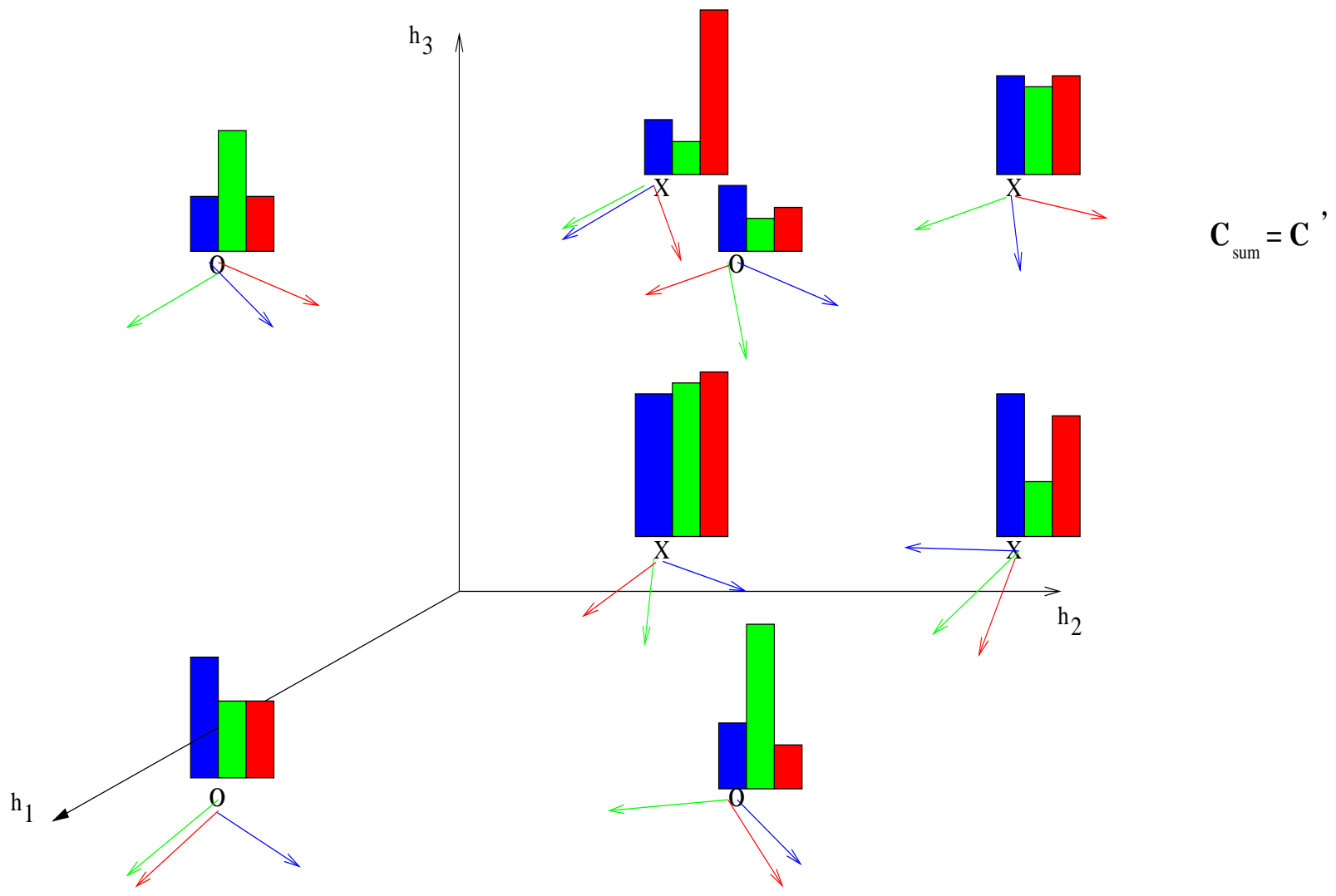
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## Joint Power and Waveform Allocation – $K \leq N$

- Optimal waveforms constitute an **orthogonal set** for any power allocation.
- Problem reduces to  $K$  independent single user [Goldsmith-Varaiya] problems, i.e.,

$$\begin{aligned} \max_{\mathbf{p}(\mathbf{h})} \quad & E_{\mathbf{h}} \left[ \sum_{i=1}^K \log \left( 1 + \frac{p_i(\mathbf{h})h_i}{\sigma^2} \right) \right] \\ \text{s.t.} \quad & E_{\mathbf{h}} [p_i(\mathbf{h})] = \bar{p}_i, \quad i = 1, \dots, K \end{aligned}$$

- Concave maximization over an affine set of constraints, using KKT conditions,

$$p_i^*(\mathbf{h}) = \left( \frac{1}{\lambda_i} - \frac{\sigma^2}{h_i} \right)^+, \quad i = 1, \dots, K$$

- Channel non-adaptive waveform selection is as good as any channel adaptive selection.

## Joint Power and Waveform Allocation – $K > N$

- For a given power control policy  $P(\mathbf{h})$ , let  $L(\mathbf{h})$  and  $\bar{L}(\mathbf{h})$  be sets of oversized and non-oversized users respectively, for a given  $\mathbf{h}$ .
- Define  $\mathbf{D} \triangleq \text{diag}(p_1 h_1, \dots, p_K h_K)$ . Optimum waveforms satisfy,

$$\mathbf{SDS}^\top \mathbf{s}_i(\mathbf{h}) = \mu_i(\mathbf{h}) \mathbf{s}_i(\mathbf{h})$$

$$\mu_i(\mathbf{h}) = \begin{cases} \frac{\sum_{j \in \bar{L}(\mathbf{h})} p_j h_j}{N - |L(\mathbf{h})|}, & i \in \bar{L}(\mathbf{h}) \\ p_i h_i, & i \in L(\mathbf{h}) \end{cases}$$

- The waveform-optimized ergodic sum-capacity is then

$$E_{\mathbf{h}} \left[ \sum_{i \in L(\mathbf{h})} \log \left( 1 + \frac{p_i(\mathbf{h}) h_i}{\sigma^2} \right) + (N - |L(\mathbf{h})|) \log \left( 1 + \frac{\sum_{i \in \bar{L}(\mathbf{h})} p_i(\mathbf{h}) h_i}{\sigma^2 (N - |L(\mathbf{h})|)} \right) \right]$$

## Maximum Number of Simultaneously Transmitting Users

**Theorem 1** *Let  $\bar{K}(\mathbf{h})$  be a subset of  $\{1, \dots, K\}$ , such that  $\forall i \in \bar{K}(\mathbf{h}), p_i^*(\mathbf{h}) > 0$ , where  $\mathbf{p}^*(\mathbf{h})$  is the maximizer of  $E_{\mathbf{h}} [C_{\text{opt}}(\mathbf{h}, \mathbf{p}(\mathbf{h}))]$ . Then, with probability 1,  $|\bar{K}(\mathbf{h})| \leq N$ .*

**Proof:**

- $C_{\text{opt}}(\mathbf{h}, \mathbf{p}(\mathbf{h}))$  is concave [Viswanath-Anantharam]
- Power constraint set is convex (affine).
- $\mathbf{p}^*(\mathbf{h})$  achieves the global optimum of the sum-capacity  $\Leftrightarrow$  it satisfies the KKT conditions.

$$\frac{h_i}{\mu_i(\mathbf{h}) + \sigma^2} \leq \lambda_i, \quad \forall \mathbf{h} \quad \text{w.e. if } p_i(\mathbf{h}) > 0$$

- Let  $|\bar{K}(\mathbf{h})| > N$ . Then, at least  $|\bar{K}(\mathbf{h})| - N + 1$  users have the same eigenvalue  $\mu_i(\mathbf{h})$ .
- Then,  $h_i/\lambda_i = h_j/\lambda_j$  for  $i \neq j, i, j \in \bar{K}(\mathbf{h})$  for at least  $|\bar{K}(\mathbf{h})| - N + 1$  users.
- This event has zero probability, therefore, with probability one,  $|\bar{K}(\mathbf{h})| \leq N$ .

## Jointly Optimum Waveforms and Powers – $K > N$

- At most  $N$  users transmit: assign **orthogonal waveforms** to those users.
- Optimum power allocation is similar to single user waterfilling

$$p_i^*(\mathbf{h}) = \begin{cases} \left( \frac{1}{\lambda_i} - \frac{\sigma^2}{h_i} \right), & i \in \bar{K}(\mathbf{h}) \\ 0, & \text{otherwise} \end{cases}$$

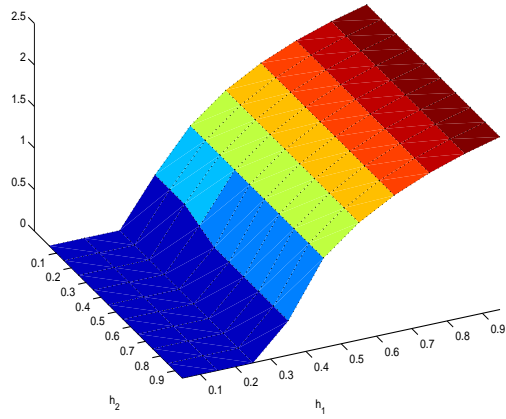
- Here, a channel adaptive allocation of orthogonal waveforms is necessary.
- Define  $\gamma_i = h_i/\lambda_i$ , and let  $\{\gamma_{[i]}\}_{i=1}^K$  be the order statistics for  $\gamma_i$ s, and let for given  $\mathbf{h}$

$$\gamma_{[1]} \geq \cdots \geq \gamma_{[n]} > \sigma^2 \geq \gamma_{[n+1]} \geq \cdots \geq \gamma_{[K+1]} = 0$$

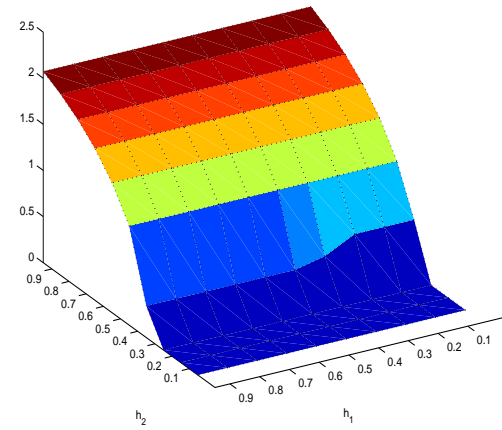
- If  $n \leq N$ , the users with highest  $n$   $\gamma_i$ 's transmit with powers  $p_i^*(\mathbf{h})$ .
- If  $n > N$ , by Theorem 1, the users with highest  $N$   $\gamma_i$ 's transmit with positive powers.

# Optimum Power Allocation: $K = 4, N = 3$

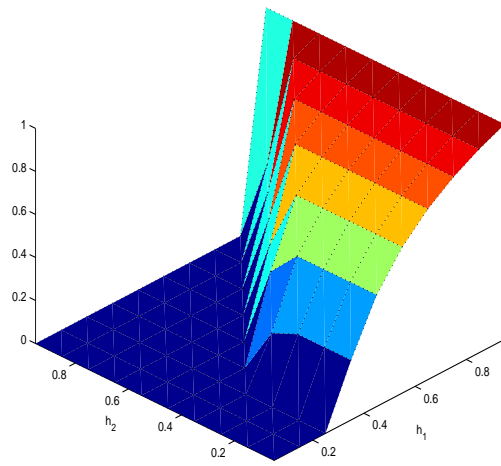
Power Allocation for User 1,  $h_3=h_4=0.4$



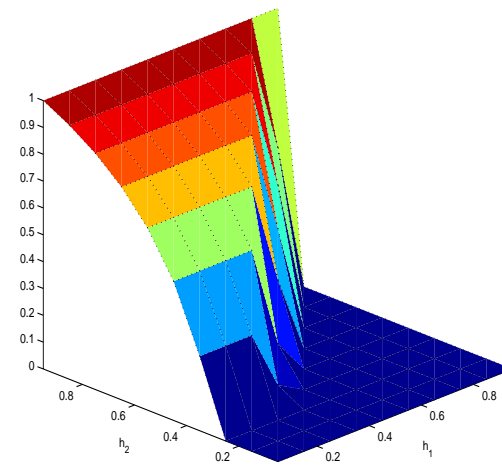
Power Allocation for User 2



Power Allocation for User 1,  $h_3=h_4=0.9$



Power Allocation for User 2,  $h_3=h_4=0.9$



## Iterative Power and Waveform Optimization

- Already characterized a “closed form” solution for optimal powers and waveforms.
- The optimum resource allocation still depends on  $\lambda_i, i = 1, \dots, K$ .
- Instead of simultaneously solving for all powers, we propose the following algorithm:

*repeat*

*for  $i = 1$  to  $K$  and for all  $\mathbf{h}$*

*-find oversized users*

*-compute waveforms for all users*

*-update  $i$ th user's power using waterfilling keeping other powers fixed*

*end*

*until  $\mathbf{p}(\mathbf{h})$  converges.*

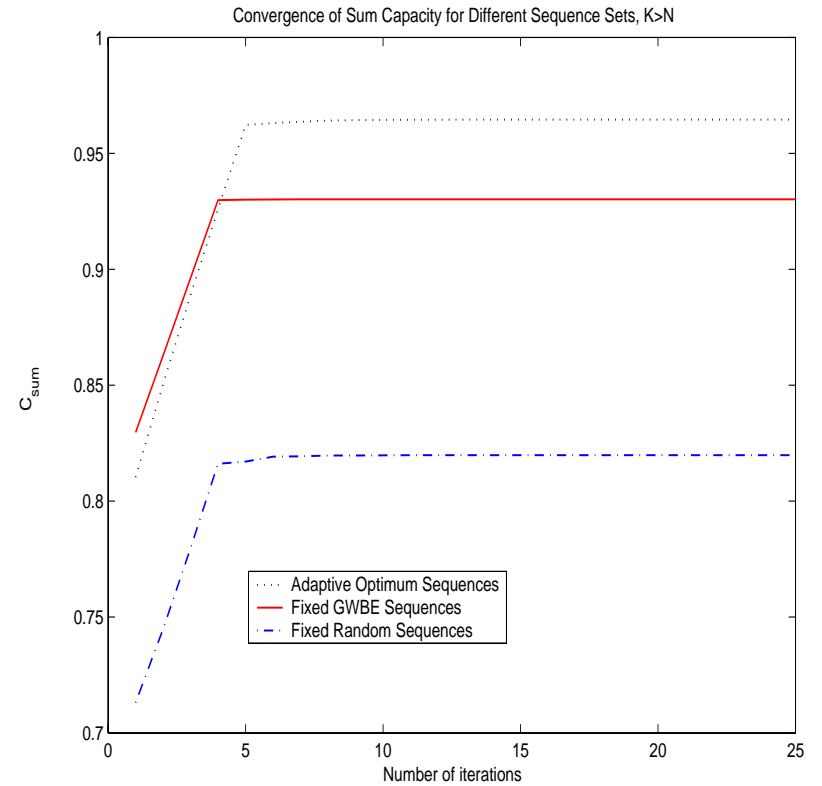
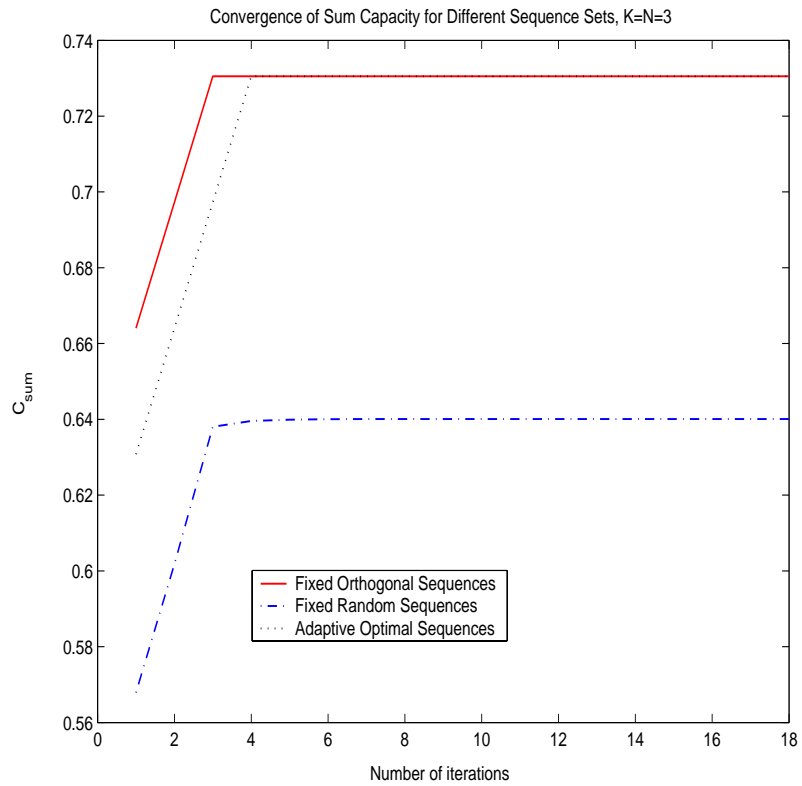
## Convergence of the Iterative Algorithm

- This algorithm corresponds to iteration of the best waveform-only update for all users and best power-only update for one user, so sum capacity values obtained are non-decreasing.
- The sum capacity is also bounded from above, so this algorithm converges to a limit.
- Same algorithm can be seen as an iterative update directly from powers-to-powers

$$p_k^{n+1}(\mathbf{h}) = \left( \frac{1}{\lambda_k} - \frac{\sigma^2 + \mu_k^n(\mathbf{h}) - h_k p_k^n(\mathbf{h})}{h_k} \right)^+$$

- The fixed point  $\mathbf{p}^{n+1}(\mathbf{h}) = \mathbf{p}^n(\mathbf{h})$  satisfies the KKT conditions for the optimization problem.
- Algorithm converges to the jointly optimum power and waveform allocation.
- **Remark:** Optimum power allocation is unique, optimum waveform allocation is not.

# Convergence and Comparison to Non-Adaptive Policies





## Summary

- Characterized optimum power allocation in fading waveform channels
  - Developed an iterative waterfilling algorithm; proved its convergence to global optimum
  - All users transmit simul. with non-zero prob. iff  $\{\mathbf{s}_i \mathbf{s}_i^\top\}_{i=1}^K$  are linearly independent
    - \*  $K \leq N$ , signatures independent: all users transmit simultaneously with  $> 0$  probability.
    - \* Maximum number of users that can transmit simul. is  $M(M + 1)/2$ ;  $M = \text{rank}(\mathbf{S})$ .
- Characterized jointly optimum power and waveform adaptation policy
  - **Optimal policy dictates orthogonal transmissions**, achieved by
    - \* **time division across fading states** [Knopp-Humblet-like]
    - \* **orthogonal waveforms for multiple users transmitting at a given state**
  - Developed an iterative algorithm; proved its convergence to global optimum
- The results may be interpreted as
  - Opportunistic scheduling in waveform channels
  - Cross-layer design: interacting/cooperating physical and MAC layers